

# Quadratic residues

David M. Burton  
- Elementary Number Theory

Def'n: (Valid for composite  $n$  also).

Let  $a \in \mathbb{Z}_p^*$  where  $p$  is an odd prime and  $a \not\equiv 0 \pmod{p}$ .

$a$  is said to be a quadratic residue modulo  $p$ , or a square modulo  $p$ , if  $\exists$  an  $x \in \mathbb{Z}_p^*$  s.t.  $x^2 \equiv a \pmod{p}$ . If no such  $x$  exists, then  $a$  is called a quadratic non-residue modulo  $p$ .  
 The set of all quadratic residues modulo  $p$  is denoted by  $\mathcal{Q}_p$  and the set of all quadratic non-residues is denoted by  $\mathcal{N}_p$ .

Note that  $0 \notin \mathbb{Z}_p^*$ , whence  $0 \notin \mathcal{Q}_p \cup \mathcal{N}_p$ .

Two basic problems dominate the theory of quadratic residues:

- i) Given a prime  $p$ , determine which  $a$  are quadratic residues mod  $p$  & which are quadratic non-residues mod  $p$ .
- ii) Given  $a$ , determine those primes  $p$  for which  $a$  is a quadratic residue mod  $p$  and those for which  $a$  is a quadratic non-residue mod  $p$ .

Let us begin  $\textcircled{i}$  with some methods for solving problem (i).

Example:

(2)

To find the quadratic residues modulo 11 we square the numbers 1, 2, ..., 10 and reduce mod 11. We obtain

$$1^2 \equiv 1, \quad 2^2 \equiv 4, \quad 3^2 \equiv 9, \quad 4^2 \equiv 5, \quad 5^2 \equiv 3 \pmod{11}$$

It suffices to square only the first half of the numbers since

$$6^2 \equiv (-5)^2 \equiv 3, \quad 7^2 \equiv (-4)^2 \equiv 5, \quad 8^2 \equiv (-3)^2 \equiv 9$$

$$9^2 \equiv (-2)^2 \equiv 4, \quad 10^2 \equiv (-1)^2 \equiv 1 \pmod{11}$$

Consequently, the quadratic residues mod 11 are 1, 3, 4, 5, 9 and the non-residues are 2, 6, 7, 8, 10.

This example illustrates the following theorem:

Th. Let  $p$  be an odd prime. Then every residue system mod  $p$  contains exactly  $(p-1)/2$  quadratic residues and exactly  $(p-1)/2$  quadratic non-residues mod  $p$ . The quadratic residues belong to the residue classes containing the numbers

$$1^2, 2^2, 3^2, \dots, \left(\frac{p-1}{2}\right)^2. \quad (1)$$

Proof. First note that the numbers in (1) are distinct mod  $p$ . In fact, if  $x^2 \equiv y^2 \pmod{p}$

With  $1 \leq x \leq \frac{p-1}{2}$  and  $1 \leq y \leq \frac{p-1}{2}$ , then

$$(x-y)(x+y) \equiv 0 \pmod{p}.$$

But  $1 < x+y < p$  so  $x-y \equiv 0 \pmod{p}$ ,  
hence  $x = y$ .

Since  $(p-k)^2 \equiv k^2 \pmod{p}$ ,

every quadratic residue is congruent mod  $p$  to  
exactly one of the numbers in (1).

This completes the proof.  $\square$   
brief

The following table of quadratic residues  $\mathbb{Q}_p$  and  
non-residues  $\bar{\mathbb{Q}}_p$  was obtained with the help of the  
above theorem.

	$p=3$	$p=5$	$p=7$	$p=11$	$p=13$
$\mathbb{Q}_p$	1	1, 4	1, 2, 4	1, 3, 4, 5, 9	1, 3, 4, 9, 10, 12
$\bar{\mathbb{Q}}_p$	2	2, 3	3, 5, 6	2, 6, 7, 8, 10	2, 5, 6, 7, 8, 11

## ④ The Legendre and Jacobi symbols

The Legendre symbol is a useful tool for keeping track of whether or not an integer  $a$  is a quadratic residue modulo a prime  $p$ .

Def. Let  $p$  be an odd prime and  $a$  an integer.

The Legendre symbol  $\left(\frac{a}{p}\right)$  is defined to be

$$\left(\frac{a}{p}\right) = \begin{cases} 0, & \text{if } p \mid a \\ 1, & \text{if } a \in \mathbb{Q}_p \\ -1, & \text{if } a \in \bar{\mathbb{Q}}_p \end{cases}$$

Examples:

$$\left(\frac{1}{p}\right) = 1, \quad \left(\frac{m^2}{p}\right) = 1, \quad \left(\frac{t}{1}\right) = -1$$

$$\left(\frac{2^2}{11}\right) = 0$$

Properties of Legendre symbol

In this we use the notation  $\chi$  for the Legendre symbol i.e.  $\chi(a) = \left(\frac{a}{p}\right)$

Let  $p$  be an odd prime and  $a, b \in \mathbb{Z}$ . Then the Legendre symbol has the following properties:

$$i) \quad \left(\frac{a}{p}\right) \equiv a^{(p-1)/2} \pmod{p}.$$

In particular,  $\left(\frac{1}{p}\right) = 1$  and  $\left(\frac{-1}{p}\right) = (-1)^{(p-1)/2}$ .

Hence  $-1 \in \mathbb{Q}_p$  if  $p \equiv 1 \pmod{4}$ , and

$-1 \in \bar{\mathbb{Q}}_p$  if  $p \equiv 3 \pmod{4}$ .

$$ii) \quad \left(\frac{ab}{p}\right) = \left(\frac{a}{p}\right) \left(\frac{b}{p}\right). \quad \text{Hence if } a \in \mathbb{Z}_p^*,$$

$$\text{then } \left(\frac{a^2}{p}\right) = 1$$

$$iii) \quad \text{If } a \equiv b \pmod{p}, \text{ then } \left(\frac{a}{p}\right) = \left(\frac{b}{p}\right)$$

$$(iv) \left(\frac{2}{p}\right) = (-1)^{\frac{(p^2-1)}{8}}$$

(5) Hence  $\left(\frac{2}{p}\right) = 1$  if  $p \not\equiv 0 \pmod{8}$

$p \equiv 1 \text{ or } 7 \pmod{8}$ , and  $\left(\frac{2}{p}\right) = -1$  if  
 $p \equiv 3 \text{ or } 5 \pmod{8}$ .

(v) (law of quadratic reciprocity)  $\rightarrow$  ~~proof~~

If  $q$  is an odd prime distinct from  $p$ , then

$$\left(\frac{p}{q}\right) = \left(\frac{q}{p}\right) (-1)^{\frac{(p-1)(q-1)}{4}}$$

In other words,  $\left(\frac{p}{q}\right) = \left(\frac{q}{p}\right)$  unless both  $p$  and  $q$  are congruent to 3 modulo 4, in which case  $\left(\frac{p}{q}\right) = -\left(\frac{q}{p}\right)$ .

The Jacobi symbol is a generalization of the Legendre symbol to integers  $n$  which are odd but not necessarily prime.

Defn: Let  $n > 3$  be odd  $\&$  with prime factorization  $n = p_1^{e_1} p_2^{e_2} \dots p_k^{e_k}$ . Then the Jacobi symbol

$$n = p_1^{e_1} p_2^{e_2} \dots p_k^{e_k}$$

$\left(\frac{a}{n}\right)$  is defined to be

$$\left(\frac{a}{n}\right) = \left(\frac{a}{p_1}\right)^{e_1} \left(\frac{a}{p_2}\right)^{e_2} \dots \left(\frac{a}{p_k}\right)^{e_k}.$$

Observe that if  $n$  is prime, then the Jacobi symbol is just the Legendre symbol.

(6)

## Properties of Jacobi Symbol.

Let  $m, n \geq 3$  be odd integers, and  $a, b \in \mathbb{Z}$ . Then the Jacobi symbol has the following properties:

- (i)  $\left(\frac{a}{n}\right) = 0, 1, \text{ or } -1$ . Moreover  $\left(\frac{a}{n}\right) = 0$  iff  $\gcd(a, n) \neq 1$ .
- (ii)  $\left(\frac{ab}{n}\right) = \left(\frac{a}{n}\right)\left(\frac{b}{n}\right)$ . Hence if  $a \in \mathbb{Z}_n^*$ , then  $\left(\frac{a^2}{n}\right) = 1$ .
- (iii)  $\left(\frac{a}{mn}\right) = \left(\frac{a}{m}\right)\left(\frac{a}{n}\right)$ .
- (iv) If  $a \equiv b \pmod{n}$ , then  $\left(\frac{a}{n}\right) = \left(\frac{b}{n}\right)$ .
- (v)  $\left(\frac{1}{n}\right) = 1$
- (vi)  $\left(\frac{-1}{n}\right) = (-1)^{\frac{(n-1)}{2}}$ . Hence  $\left(\frac{-1}{n}\right) = 1$  if  $n \equiv 1 \pmod{4}$ , and  $\left(\frac{-1}{n}\right) = -1$  if  $n \equiv 3 \pmod{4}$ .
- (vii)  $\left(\frac{2}{n}\right) = (-1)^{\frac{(n^2-1)/8}{2}}$ . Hence  $\left(\frac{2}{n}\right) = 1$  if  $n \equiv 1 \text{ or } 7 \pmod{8}$ , and  $\left(\frac{2}{n}\right) = -1$  if  $n \equiv 3 \text{ or } 5 \pmod{8}$ .
- (viii)  $\left(\frac{m}{n}\right) = \left(\frac{n}{m}\right)(-1)^{\frac{(m-1)(n-1)/4}{2}}$ . In other words,  $\left(\frac{m}{n}\right) = \left(\frac{n}{m}\right)$  unless both  $m$  &  $n$  are congruent to 3 modulo 4, in which case  $\left(\frac{m}{n}\right) = -\left(\frac{n}{m}\right)$ .

(7)

It is clear that  $\left(\frac{a}{p}\right) = \left(\frac{b}{p}\right)$  whenever  $a \equiv b \pmod{p}$ ,  
 So  $\left(\frac{a}{p}\right)$  is a periodic function of  $a$  with period  $p$ .

Fermat's little theorem (A special case of Euler's Th.)

$$a^{p-1} \equiv 1 \pmod{p} \text{ if } p \nmid a.$$

$$\text{Since } a^{p-1} - 1 = (a^{\frac{p-1}{2}} + 1)(a^{\frac{p-1}{2}} - 1)$$

$$\text{it follows that } a^{\frac{p-1}{2}} \equiv \pm 1 \pmod{p}.$$

The next theorem tells us that we get  $+1$  if  $a \in \mathbb{Q}_p$   
 and  $-1$  if  $a \in \bar{\mathbb{Q}}_p$ .

Theorem (Euler's criterion)

Let  $p$  be an odd prime. Then for all  $a$  we have

$$\left(\frac{a}{p}\right) \equiv a^{\frac{p-1}{2}} \pmod{p}.$$

proof. Case 1 If  $a \equiv 0 \pmod{p}$  the result is trivial  
 since both members are congruent to 0 mod  $p$ .

Case 2 Now suppose that  $\left(\frac{a}{p}\right) = 1$ .

Then  $\exists x$  s.t.  $x^2 \equiv a \pmod{p}$  and hence

$$a^{\frac{p-1}{2}} \equiv (x^2)^{\frac{p-1}{2}} = x^{p-1} \equiv 1 \pmod{p}$$

by Fermat's little theorem.

Thus the theorem ~~is~~ is proved if  $\left(\frac{a}{p}\right) = 1$

Case 3 Now suppose that  $\left(\frac{a}{p}\right) = -1$  & consider the  
 following polynomial

⑧

$$f(x) = x^{\frac{p-1}{2}} - 1$$

Since  $f(x)$  is a polynomial of degree  $\frac{p-1}{2}$ , the congruence

$f(x) \equiv 0 \pmod{p}$  (by Lagrange's Th.) (not true for composite numbers)  
 has at most  $\frac{p-1}{2}$  solutions. But the  $\frac{p-1}{2}$  quadratic residues mod  $p$  are solutions, so the non-residues are not.

Hence  $a^{\frac{p-1}{2}} \not\equiv 1 \pmod{p}$ .

But  $a^{\frac{p-1}{2}} \equiv \pm 1 \pmod{p}$  so  $a^{\frac{p-1}{2}} \equiv -1 \pmod{p}$

This completes the proof.

Theorem Legendre's symbol  $\left(\frac{a}{p}\right)$  is a completely multiplicative function of  $a$ . [i.e.  $\left(\frac{ab}{p}\right) = \left(\frac{a}{p}\right)\left(\frac{b}{p}\right)$ ].

Proof Case 1

If  $p|a$  or  $p|b$ , then  $p|ab$

so  $\left(\frac{ab}{p}\right) = 0$  and either  $\left(\frac{a}{p}\right) = 0$  or  $\left(\frac{b}{p}\right) = 0$ .

Therefore,  $\left(\frac{ab}{p}\right) = \left(\frac{a}{p}\right)\left(\frac{b}{p}\right)$  if  $p|a$  or  $p|b$ .

Case 2 If  $p \nmid a$  and  $p \nmid b$ , then  $p \nmid ab$  we have

$$\left(\frac{ab}{p}\right) = (ab)^{\frac{p-1}{2}} = a^{\frac{p-1}{2}} b^{\frac{p-1}{2}} \equiv \left(\frac{a}{p}\right)\left(\frac{b}{p}\right) \pmod{p}$$

But each of  $\left(\frac{ab}{p}\right)$ ,  $\left(\frac{a}{p}\right)$  &  $\left(\frac{b}{p}\right)$  is 1 or -1.  
 So the difference

$$\left(\frac{ab}{p}\right) - \left(\frac{a}{p}\right)\left(\frac{b}{p}\right)$$

is either 0, 2 or -2.  
 Since this difference is divisible by  $p$ , it must be 0.

## Evaluation of $\left(\frac{-1}{p}\right)$ and $\left(\frac{2}{p}\right)$

(9)

Theorem for every odd prime  $p$  we have

$$\left(\frac{-1}{p}\right) = (-1)^{\frac{p-1}{2}} = \begin{cases} 1 & \text{if } p \equiv 1 \pmod{4} \\ -1 & \text{if } p \equiv 3 \pmod{4}. \end{cases}$$

Proof. By Euler's criterion we have

$$\left(\frac{-1}{p}\right) \equiv (-1)^{\frac{p-1}{2}} \pmod{p}.$$

Since each member of this congruence is  $1$  or  $-1$ , the two members are equal.

Theorem for every odd prime  $p$  we have

$$\left(\frac{2}{p}\right) = (-1)^{\frac{p^2-1}{8}} = \begin{cases} 1 & \text{if } p \equiv \pm 1 \pmod{8} \\ -1 & \text{if } p \equiv \pm 3 \pmod{8}. \end{cases}$$

Proof. Consider the following  $\frac{p-1}{2}$  congruences:

$$p-1 \equiv 1 \pmod{p}$$

$$2 \equiv 2 \pmod{p}$$

$$p-3 \equiv 3 \pmod{p}$$

$$4 \equiv 4 \pmod{p}$$

$$\gamma \equiv \frac{p-1}{2} (-1)^{\frac{p-1}{2}} \pmod{p}$$

where  $\gamma$  is either  $p - \frac{p-1}{2}$  or  $\frac{p-1}{2}$ .

Multiplying these together & note that each integer on the left is even. We obtain

$$2 \cdot 4 \cdot 6 \cdots (p-1) \equiv \left\lfloor \frac{p-1}{2} \right\rfloor (-1)^{1+2+\cdots+\frac{p-1}{2}} \pmod{p}$$

This gives us

$$2^{\frac{p-1}{2}} \left\lfloor \frac{p-1}{2} \right\rfloor = \left\lfloor \frac{p-1}{2} \right\rfloor (-1)^{\frac{p-1}{8}} \pmod{p}.$$

Since  $\left\lfloor \frac{p-1}{2} \right\rfloor \neq 0 \pmod{p}$ , this implies

$$2^{\frac{p-1}{2}} \equiv (-1)^{\frac{p-1}{8}} \pmod{p}.$$

By Euler's criterion, we have

$$2^{\frac{p-1}{2}} \equiv \left(\frac{2}{p}\right) \pmod{p},$$

and since each member is 1 or -1, the two members are equal. This completes the proof.

(11)

## Jacobi Symbol

Defn:Let  $n > 3$  be odd with prime factorization

$$n = p_1^{e_1} p_2^{e_2} \cdots p_k^{e_k}.$$

Then the Jacobi symbol  $\left(\frac{a}{n}\right)$  is defined to be

$$\left(\frac{a}{n}\right) = \left(\frac{a}{p_1}\right)^{e_1} \left(\frac{a}{p_2}\right)^{e_2} \cdots \left(\frac{a}{p_k}\right)^{e_k}.$$

Case  $n$  prime  $\rightarrow$  Jacobi symbol is just the Legendre symbol.

## Fact (Properties of Jacobi symbol)

Let  $m > 3, n > 3$  be odd integers, and  $a, b \in \mathbb{Z}$ .

Then the Jacobi symbol has the following properties:

(i)  $\left(\frac{a}{n}\right) = 0, 1 \text{ or } -1$ .

$$\left(\frac{a}{n}\right) = 0 \text{ iff } \gcd(a, n) \neq 1.$$

(ii)  $\left(\frac{ab}{n}\right) = \left(\frac{a}{n}\right) \left(\frac{b}{n}\right)$  [Hence if  $a \in \mathbb{Z}_n^*$ , then  $\left(\frac{a^2}{n}\right) = 1$ ]

(iii)  $\left(\frac{a}{mn}\right) = \left(\frac{a}{m}\right) \left(\frac{a}{n}\right)$ .

(iv) If  $a \equiv b \pmod{n}$ , then  $\left(\frac{a}{n}\right) = \left(\frac{b}{n}\right)$ .

(v)  $\left(\frac{1}{n}\right) = 1$

(vi)  $\left(-\frac{1}{n}\right) = (-1)^{\frac{n-1}{2}}$  [Hence  $\left(-\frac{1}{n}\right) = \begin{cases} 1 & \text{if } n \equiv 1 \pmod{4} \\ -1 & \text{if } n \equiv 3 \pmod{4} \end{cases}$ ]

(vii)  $\left(\frac{2}{n}\right) = (-1)^{\frac{n^2-1}{8}}$ . [Hence (12)  $\left(\frac{2}{n}\right) = \begin{cases} 1 & \text{if } n \equiv 1 \pmod{8} \\ -1 & \text{if } n \equiv 3 \pmod{8} \text{ or } 5 \end{cases}$ ]

(viii)  $\left(\frac{m}{n}\right) \left(\frac{n}{m}\right) = (-1)^{\frac{(m-1)(n-1)}{4}}$

i.e.  $\left(\frac{mn}{n}\right) = \left(\frac{n}{mn}\right)$  unless both  $m, n$  are congruent to 3 modulo 4, in which case  $\left(\frac{mn}{n}\right) = -\left(\frac{n}{m}\right)$ .

Note: If  $n$  is odd and  $a = 2^e a_1$  when  $a_1$  is odd,

then

$$\left(\frac{a}{n}\right) = \left(\frac{2^e}{n}\right) \left(\frac{a_1}{n}\right)$$

$$= \cancel{\left(\frac{2}{n}\right)^e} \cancel{\left(\frac{a_1}{n}\right)} \\ = \left(\frac{2}{n}\right)^e \left(\frac{n \bmod a_1}{a_1}\right) (-1)^{\frac{(q-1)(m_1-1)}{4}}$$

This observation yields the following recursive alg m. for computing  $\left(\frac{a}{n}\right)$ , which does not require the prime factorization of  $n$ .

Note -

deterministic (13)

no poly-time algm. to find quadratic non-residues modulo a prime  $p$ ,

although we know that half of the elements in  $\mathbb{Z}_p^*$  are quadratic residues modulo  $p$ .

Randomized algm. for finding quadratic non-residues modulo  $p$

a  $\in \mathbb{Z}_p^* \rightarrow$  random selection of a until 1 one is found satisfying

$$\left(\frac{a}{p}\right) = -1.$$

Example -

Jacobi Symbol Computation

$$a = 158, n = 235.$$

$$\begin{aligned} \left(\frac{158}{235}\right) &= \left(\frac{2}{235}\right) \left(\frac{79}{235}\right) = (-1) \cancel{\left(\frac{79}{235}\right)} \\ &= (-1) \left(\frac{235}{79}\right) (-1) \\ &= (-1) \left(\frac{77}{79}\right) (-1) \\ &= \left(\frac{77}{79}\right) = \left(\frac{79}{77}\right) (-1) \frac{(77-1)(79-1)}{4} = \left(\frac{2}{77}\right) = -1 \end{aligned}$$

$$8 \mid \frac{77}{72} \mid 9$$

$$8 \mid \frac{235}{16} \mid 29$$

$$\begin{array}{r} 79 \mid \frac{235}{158} \mid 2 \\ \hline 77 \\ 38 \times 234 \\ \hline 127 \\ 38 \times 78 \\ \hline 1 \end{array}$$

# Algorithm      Jacobi Symbol $\leftrightarrow$ (Legendre symbol)

Computation

---

$\text{JACOBI } (a, n)$

Input: an odd integer  $n \geq 3$ , and an integer  $a$ , ~~such that~~

$$0 \leq a < n$$

Output: the Jacobi symbol  $\left(\frac{a}{n}\right)$ . (hence the Legendre symbol when  $n$  is prime).

1. If  $a=0$  then return(0)
2. If  $a=1$  then return(1)
3. Write  $a = 2^e a_1$ , where  $a_1$  is odd.
4. If  $e$  is even, then set  $\rho \leftarrow 1$ .  
otherwise, set  $\rho \leftarrow 1$  if  $n \equiv 1 \pmod{8}$ ,  
or set  $\rho \leftarrow -1$  if  $n \equiv 3 \pmod{8}$ .
5. If  $n \equiv 3 \pmod{4}$  and  $a_1 \equiv 3 \pmod{4}$  then set  $\rho \leftarrow -\rho$
6. Set  $n_1 \leftarrow n \bmod a_1$
7. If  $a_1 = 1$  then return(1)  
otherwise return( $\rho, \text{JACOBI } (a_1, n_1)$ )

Note:  $\left(\frac{a}{n}\right) = 1 \not\Rightarrow a \in \mathbb{Q}_n$       Example  
 But  $a \in \mathbb{Q}_n \Rightarrow \left(\frac{a}{n}\right) = 1$

$$\left\{ \begin{array}{l} \left(\frac{5}{21}\right) = 1 \\ \text{but } 5 \notin \mathbb{Q}_{21} \end{array} \right.$$

$$\left(\frac{5}{3}\right) \left(\frac{7}{3}\right) = \left(\frac{21}{3}\right) \left(\frac{5}{7}\right) = (-1)^2 = 1$$

Defn: Let  $n \geq 3$  be an odd integer, and let  $J_n = \{at z_n^* \mid \left(\frac{a}{z}\right) = 1\}$ .  
 The set of pseudosquares modulo  $n$ , denoted by  $\tilde{\mathbb{Q}}_n$  is defined to be the set  $J_n - \mathbb{Q}_n$ .

Fact: If  $n = pq$ ,  $p, q$  are two distinct primes.

$$\text{Then } |\tilde{\mathbb{Q}}_n| = |\mathbb{Q}_n| = \frac{(p-1)(q-1)}{4}$$

i.e. half of the elements in  $J_n$  are quadratic residues (squares) and the other half are pseudosquares modulo  $n$ .

$$\left(\frac{59}{67}\right) = \left(\frac{67}{59}\right)(-1) = \left(\frac{8}{59}\right)$$

$$= \cancel{\left(\frac{1}{59}\right)} = \left(\frac{2}{59}\right)^3 = (-1)^3 = -1$$

$$\frac{66}{2} \quad \frac{58}{2}$$

$$8 \mid 59 \mid 7 \quad \frac{56}{3}$$

$$8k+3$$

Theorem If  $P$  is an odd positive integer, then (16)

$$\left(-\frac{1}{P}\right) = (-1)^{\frac{P-1}{2}}$$

$$+ \left(\frac{2}{P}\right) = (-1)^{\frac{P^2-1}{8}}$$

Proof- Write  $P = p_1 p_2 \dots p_m$  where the prime factors  $p_i$  are not necessarily distinct.

$$P = \prod_{i=1}^m p_i$$

$$= \prod_{i=1}^m (1 + p_i - 1) = 1 + \sum_{i=1}^m (p_i - 1) + \sum_{i=1}^m \sum_{j=1, j \neq i}^m (p_i - 1)(p_j - 1) + \dots$$

Each  $(p_i - 1)$  is even.

$$\therefore P \equiv 1 + \sum_{i=1}^m (p_i - 1) \pmod{4}$$

$a \equiv b \pmod{2}$   
 $\Rightarrow a, b$  both odd  
 or both even  
 $(-1)^n = (-1)^3$ .

$$\boxed{\text{or } \frac{1}{2}(P-1) \equiv \frac{1}{2} \sum_{i=1}^m (p_i - 1) \pmod{2}} \quad \text{--- (1)}$$

$$\therefore \left(-\frac{1}{P}\right) = \prod_{i=1}^m \left(\frac{-1}{p_i}\right) = \prod_{i=1}^m (-1)^{\frac{(p_i-1)}{2}} = (-1)^{\sum_{i=1}^m \frac{(p_i-1)}{2}} = \boxed{(-1)^{\frac{P-1}{2}}}$$

or  $d = m \cdot m'$  or  $-1 \not\mid d$  then then  
if Theorem if  $ac \equiv bc \pmod{m}$  &  $\gcd(c, m) > 1$   
then  $a \equiv b \pmod{\frac{m}{d}}$

Corollary if  $ac \equiv bc \pmod{m}$  &  $c \mid m$ , then  
 $a \equiv b \pmod{\frac{m}{c}}$

$$P^2 = \prod_{i=1}^m (1 + p_i^2 - 1)$$

$$= 1 + \sum_{i=1}^m (p_i^2 - 1) + \sum_{i=1}^m \sum_{j=1, j \neq i}^m (p_i^2 - 1)(p_j^2 - 1) + \dots$$

$\therefore p_i$  is odd,

each  $(p_i^2 - 1)$  is even divisible by 8 i.e.  $p_i^2 - 1 \equiv 0 \pmod{8}$

$$\therefore P^2 \equiv 1 + \sum_{i=1}^m (p_i^2 - 1) \pmod{8}$$

~~$a \equiv b \pmod{m}$~~

$$8 \times \frac{1}{8} (P^2 - 1) \equiv 8 \times \frac{1}{8} \sum_{i=1}^m (p_i^2 - 1) \pmod{8 \times 8}$$

$$c = 8 | 64.$$

$$8 \rightarrow a - b = 8k$$

$$\Rightarrow \boxed{\frac{1}{8} (P^2 - 1) \equiv \frac{1}{8} \sum_{i=1}^m (p_i^2 - 1) \pmod{2}} - (B)$$

$$\therefore \left(\frac{2}{P}\right) = \prod_{i=1}^m \left(\frac{2}{p_i}\right) = \prod_{i=1}^m (-1)^{(p_i^2 - 1)/8} = (-1)^{\frac{1}{8} \sum (p_i^2 - 1)} = (-1)^{\frac{1}{8} (P^2 - 1)}$$

Theorem (Reciprocity Law for Jacobi Symbols)

If  $P$  and  $Q$  are ~~primes~~ positive odd integers with  $\gcd(P, Q) = 1$ , then  $\left(\frac{P}{Q}\right) \left(\frac{Q}{P}\right) = (-1)^{\frac{(P-1)(Q-1)}{4}}$ .

Proof.

Let  $P = p_1 p_2 \dots p_m$ ,  $Q = q_1 q_2 \dots q_n$ , where  $p_i, q_i$  are primes.

$$\therefore \left(\frac{P}{Q}\right) \left(\frac{Q}{P}\right) = \prod_{i=1}^m \prod_{j=1}^n \left(\frac{p_i}{q_j}\right) \left(\frac{q_j}{p_i}\right) = (-1)^r, \text{ say.}$$

$$r = \sum_{i=1}^m \sum_{j=1}^n \frac{(p_i - 1)(q_j - 1)}{4} = \sum_{i=1}^m \frac{(p_i - 1)}{2} \sum_{j=1}^n \frac{(q_j - 1)}{2} = \frac{1}{4} (P-1)(Q-1) \quad (\text{why?})$$

$\nexists r, \frac{1}{4}(P-1)(Q-1)$  have the same parity.

[by (A)]

Square of an odd int. is often  
form  $8k+1$

$p_i \rightarrow 4k, 4k+1, 4k+2, 4k+3$

Proof:  $4k, 4k+1, 4k+2, 4k+3$

odd  
sq.  
 $8k+1$

(18)

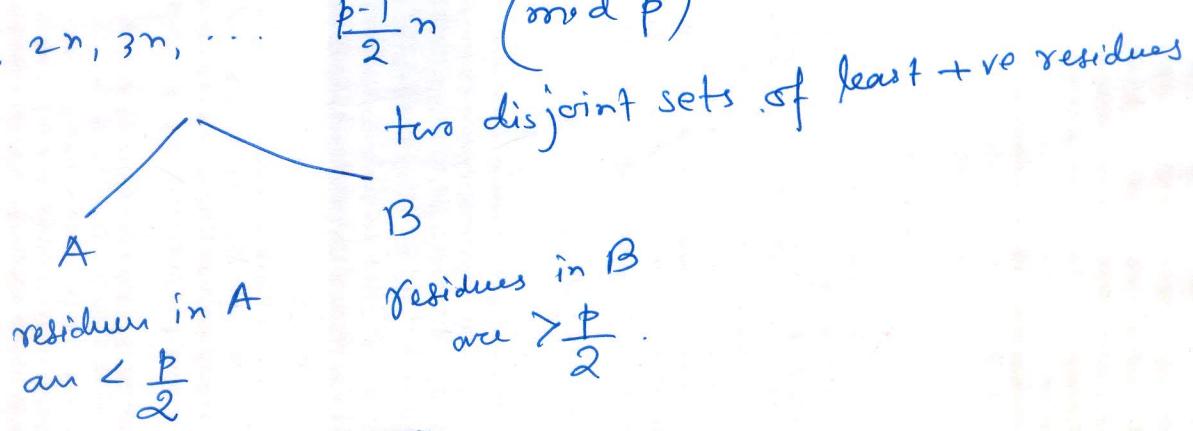
Theorem (Gauss Lemma) Let  $p$  be an odd prime.  
 Assume  $n \not\equiv 0 \pmod{p}$  and consider the least positive integer  
 mod  $p$  of the following  $(p-1)/2$  multiples of  $n$ :

$$n, 2n, 3n, \dots, \frac{p-1}{2}n.$$

If  $m$  denotes the number of these residues which exceeds  $p/2$ ,  
 then  $\left(\frac{n}{p}\right) = (-1)^m$ .

Proof - Consider least +ve residues mod  $p$  of numbers

$$n, 2n, 3n, \dots, \frac{p-1}{2}n \pmod{p}$$



Let  $A = \{a_1, a_2, \dots, a_k\}$

When  $a_i \equiv tn \pmod{p}$  for some  $t_2 \leq \frac{p-1}{2}$  and  $0 < a_i < \frac{p}{2}$ .

Let  $B = \{b_1, b_2, \dots, b_m\}$

When  $b_i \equiv tn \pmod{p}$  for some  $t_2 \leq \frac{p-1}{2}$  and  $\frac{p}{2} < b_i < p$ .

As  $A, B$  are disjoint,

$$m+k = \frac{p-1}{2}$$

Let  $C = \{c_1, c_2, \dots, c_m\}$

As  $\frac{p}{2} < b_i < p$ ,  $0 < c_i < \frac{p}{2}$

i.e. elements of  $C$  lie in the same interval as the elements of  $A$ .

$$\text{when } c_i = p - b_i$$

Claim  $A, C$  are disjoint.

(19)

If not,  $c_i = a_j$  for some pair  $i, j$ .

$$\text{Then } p - b_i = a_j$$

$$\text{or } a_j + b_i \equiv 0 \pmod{p} \quad \dots \textcircled{1}$$

$$\text{Let } a_i \equiv tn \pmod{p}, \quad 1 \leq t \leq \frac{p-1}{2} \quad 0 < t < \frac{p}{2}$$

$$b_i \equiv rn \pmod{p}, \quad 1 \leq r \leq \frac{p-1}{2} \quad 0 < r < \frac{p}{2}$$

$$\text{Then } \textcircled{1} \Rightarrow \cancel{tn + rn} = \cancel{(t+r)n}$$

$$\therefore \textcircled{1} \Rightarrow tn + rn \equiv (t+r)n \equiv 0 \pmod{p}$$

( $\rightarrow \leftarrow$ )

as  $p \nmid n$  and  $p \nmid (t+r)$

$$0 < t+r < p$$

$\therefore A, C$  are disjoint sets.

$$AVC = \left\{ \underbrace{a_1, a_2, \dots, a_k}_{\substack{0 < a_i < \frac{p}{2} \\ i.e. 1 \leq a_i \leq \frac{p-1}{2}}}, \underbrace{c_1, c_2, \dots, c_m}_{\substack{0 < c_i < \frac{p}{2} \\ i.e. 1 \leq c_i \leq \frac{p-1}{2}}} \right\} = \left\{ 1, 2, \dots, \frac{p-1}{2} \right\}$$

$$|AVC| = m+k = \frac{p-1}{2}$$

$$\therefore a_1 a_2 \dots a_k c_1 c_2 \dots c_m = 1 \cdot 2 \cdot \dots \cdot \frac{p-1}{2}$$

$$a_1 a_2 \dots a_k (p-b_1) (p-b_2) \dots (p-b_m) = \underline{\frac{p-1}{2}}$$

$$\therefore \underline{\frac{p-1}{2}} \equiv a_1 a_2 \dots a_k (-b_1) (-b_2) \dots (-b_m) \pmod{p}$$

$$\equiv (-1)^m a_1 a_2 \dots a_k b_1 b_2 \dots b_m \pmod{p}$$

$$\text{So } \frac{p-1}{2} \equiv (-1)^m \pmod{p}$$

$$\frac{\binom{n}{p}}{\binom{m}{p}} \text{ by Euler's criterion} \equiv (-1)^m (n)(2^n)(3^n) \dots \left(\frac{p-1}{2} n\right) \pmod{p}$$

$$\equiv (-1)^m n^{\frac{p-1}{2}} \underline{\frac{p-1}{2}} \pmod{p} \cdot \frac{\binom{n}{p}}{\binom{m}{p}} = (-1)^m$$

### Theorem

exact value of  $m$  is not needed, only its parity is required.

(20) Let  $m$  be the number defined in Gram's Lemma. Then

$$m = \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{tn}{p} \right] + (n-1) \frac{\frac{p-1}{2}}{8} \pmod{2}.$$

In particular, if  $n$  is odd we have

$$m \equiv \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{tn}{p} \right] \pmod{2}.$$

Proof- Note that  $m$  is the no. of least +ve integer residues mod  $p$

the numbers

$$n, 2n, 3n, \dots, \frac{(p-1)}{2}n.$$

which exceeds  $\frac{p}{2}$ .

$$\frac{tn}{p} = \left[ \frac{tn}{p} \right] + \left\{ \frac{tn}{p} \right\}, \quad 0 < \left\{ \frac{tn}{p} \right\} < 1.$$

↓  
integral part      ↓  
fraction part

$$\Rightarrow tn = p \left[ \frac{tn}{p} \right] + p \left\{ \frac{tn}{p} \right\} = p \left[ \frac{tn}{p} \right] + r_t, \text{ say,}$$

where  $0 \leq r_t < p$ .

$\therefore r_t = tn - p \left[ \frac{tn}{p} \right] \rightarrow$  least +ve residue of  $tn$  mod  $p$ .

$$\therefore \left\{ r_1, r_2, \dots, r_{\frac{p-1}{2}} \right\} = \{a_1, a_2, \dots, a_k; b_1, b_2, \dots, b_m\}$$

$$\text{Also } \left\{ 1, 2, \dots, \frac{p-1}{2} \right\} = \{a_1, a_2, \dots, a_k; c_1, c_2, \dots, c_n\}$$

whence  $c_i = p - b_i$ .

$$\therefore \sum_{t=1}^{\frac{p-1}{2}} r_t = \sum_{i=1}^k a_i + \sum_{j=1}^m b_j$$

(2)

$$\ell \sum_{t=1}^{\frac{p-1}{2}} t = \sum_{i=1}^k a_i + \sum_{j=1}^m c_j = \sum_{i=1}^k a_i + mp - \sum_{j=1}^m b_j$$

or we get

$$\sum_{i=1}^k a_i + \sum_{j=1}^m b_j = \sum_{t=1}^{\frac{p-1}{2}} \left( tn - p \left[ \frac{tn}{p} \right] \right) = n \sum_{t=1}^{\frac{p-1}{2}} t - p \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{tn}{p} \right]$$

$$mp + \sum_{i=1}^k a_i - \sum_{j=1}^m b_j = \sum_{t=1}^{\frac{p-1}{2}} t$$

~~or adding we get~~

or adding we get

~~$$-mp + \sum_{i=1}^k a_i = n+1$$~~

$$mp + \sum_{i=1}^k a_i = (n+1) \sum_{t=1}^{\frac{p-1}{2}} t - p \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{tn}{p} \right]$$

$$= (n+1) \frac{\frac{p-1}{2}}{8} - p \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{tn}{p} \right]$$

$$\text{mod } 3 \quad n+1 \equiv n-1 \pmod{2} \quad \ell \quad p \equiv 1 \pmod{2}$$

$$\frac{p-1}{2}$$

$$\therefore m \equiv (n-1) \frac{\frac{p-1}{2}}{8} + p \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{tn}{p} \right] \pmod{3}$$

### Theorem (Quadratic Reciprocity Law)

If  $p \neq q$  are distinct odd primes, then

$$\left(\frac{p}{q}\right) \left(\frac{q}{p}\right) = (-1)^{\frac{(p-1)(q-1)}{4}}.$$

(22)

Proof- By Gauss Lemma,  $\left(\frac{q}{p}\right) = (-1)^m$

$$\text{where } m \equiv \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{+q}{p} \right] \pmod{2}.$$

$$\text{Similarly, } \left(\frac{p}{q}\right) = (-1)^n$$

$$\text{where } n \equiv \sum_{s=1}^{\frac{q-1}{2}} \left[ \frac{+p}{q} \right] \pmod{2}.$$

$\therefore \left(\frac{p}{q}\right) \left(\frac{q}{p}\right) = (-1)^{m+n}$  and the proof follows if the following claim is proved.

$$\text{Claim: } \sum_{t=1}^{\frac{p-1}{2}} \left[ \frac{+q}{p} \right] + \sum_{s=1}^{\frac{q-1}{2}} \left[ \frac{+p}{q} \right] = \frac{p-1}{2} \frac{q-1}{2}$$

$$\text{Let } f(x, y) = qy - px$$

If  $x, y$  are non-zero integers, then  $f(x, y)$  is a non-zero integer.

As  $x$  takes values  $1, 2, \dots, \frac{p-1}{2}$   $y$  takes values  $1, 2, \dots, \frac{q-1}{2}$ ,

$f(x, y)$  takes  $\frac{p-1}{2} \cdot \frac{q-1}{2}$  values no two of which are equal as  $f(x_1, y_1) - f(x_2, y_2) = f(x_1 - x_2, y_1 - y_2) \neq 0$  for  $x_1 \neq x_2$  and  $y_1 \neq y_2$ .

$f(x,y) \xrightarrow{+P}$  total  $\frac{p-1}{2} \cdot \frac{q-1}{2}$  distinct values  
 as  $x \rightarrow 1, \dots, \frac{p-1}{2}$   
 and  $y \rightarrow 1, \dots, \frac{q-1}{2}$ .  
 how many are  $< 0$  & how many are  $> 0$ .

(2)

for each  $\Rightarrow$  fixed  $x$ ,

$$f(x,y) > 0 \text{ iff } qy < \frac{qx}{p}$$

$$\text{or } y \leq \left[ \frac{\frac{qx}{p}}{q} \right]$$

$\therefore$  Total no. of +ve values is  $\sum_{x=1}^{\frac{p-1}{2}} \left[ \frac{\frac{qx}{p}}{q} \right]$

$$\sum_{y=1}^{\frac{q-1}{2}} \left[ \frac{\frac{py}{q}}{p} \right]$$

Similarly, total no. of -ve values is

$$\therefore \frac{p-1}{2} \cdot \frac{q-1}{2} = \sum_{x=1}^{\frac{p-1}{2}} \left[ \frac{\frac{qx}{p}}{q} \right] + \sum_{y=1}^{\frac{q-1}{2}} \left[ \frac{\frac{py}{q}}{p} \right]$$

Hence proved.

## 24 Security: factoring hard

### Rabin Cryptosystem

$n = p q$ ,  $p, q$  primes,  $p, q \equiv 3 \pmod{4}$

$$\mathcal{P} = \mathcal{C} = \mathbb{Z}_n^*, \quad \mathcal{K} = \left\{ \underbrace{(n, p, q)}_{\substack{\downarrow \text{public} \\ \downarrow \text{secret}}} \right\}.$$

for  $K = (n, p, q)$ ,

$$e_K(x) = x^2 \pmod{n}$$

$$d_K(y) = \sqrt{y} \pmod{n}$$

This restriction  
simplifies some  
aspects of computer  
analysis of  
the cryptosystem.

Note: Restriction on  
 $p, q$  can be  
omitted.

$$\mathcal{P} = \mathcal{C} = \mathbb{Z}_n^*$$

### Decryption

$$y = x^2 \pmod{n}, \quad x \text{ unknown.}$$

→ quadratic eqn. in  $\mathbb{Z}_n$ .

decryption requires ~~to~~ extracting  
sq. roots modulo  $n$ .

equivalent + solve using (Chinese remainder th.)

$$\begin{cases} y = x_1^2 \pmod{p} \\ y = x_2^2 \pmod{q} \end{cases} \quad \text{i.e. if } y \text{ is a quadratic residue modulo } p \text{ and modulo } q.$$

$a \in \mathbb{Z}_n^* \iff a \pmod{n}$   
 $\exists x \in \mathbb{Z}_n^* \text{ s.t. } a = x^2 \pmod{n}$   
 $a = x^2 \pmod{n} \iff a$  is QR

CRT.  
(Chinese)  
algm.

$m_1, \dots, m_r \rightarrow$  pairwise  
relatively prime  
 $a_1, \dots, a_r \rightarrow a_i \in \mathbb{Z}$ .

$$\left. \begin{array}{l} x = a_1 \pmod{m_1} \\ x = a_2 \pmod{m_2} \\ \vdots \\ x = a_r \pmod{m_r} \end{array} \right\} \text{ has a unique soln. modulo } M = m_1 \cdots m_r$$

$$x = \sum_{i=1}^{r-1} a_i M_i y_i \pmod{M}$$

$$\text{where } M_i = M/m_i, \quad y_i = M_i^{-1} \pmod{m_i} \quad \text{for } 1 \leq i \leq r.$$

- $\phi = 3 \bmod 4$ ,  $\alpha^2 \equiv 3 \bmod 4$ .

decryption is easy.

$$\left(\pm y^{\frac{p+1}{4}}\right)^2 = y^{\frac{p+1}{2}} \stackrel{(1)}{=} y \cdot y^{\frac{p-1}{2}} = y \bmod p$$

As  $\left(\frac{y}{p}\right) \equiv y^{\frac{p-1}{2}} \equiv 1 \bmod \phi$  by

~~Fermat's Little theorem~~

Euler's criterion.

If  $y$  is a quadratic residue modulo  $p$ .

& in this case  $\pm y^{\frac{p+1}{4}}$  are the two sq. roots of  $y \bmod p$

Similarly,  $\pm y^{\frac{q+1}{4}}$  are the two sq. roots of  $y \bmod q$

• find 4 sq. roots modulo  $n = pq$  using CRT.

• Euler's criterion yields only an answer  
'yes' or 'no'  $\rightarrow$  quadratic Residue or not?

$\rightarrow$  does not help us to compute the sq. roots of ~~this~~ modulo  $n$ .

•  ~~$\phi, \alpha \neq 3 \bmod 4$~~

if  $p = 4k_1 + 3$ ,  $q = 4k_2 + 3$

$\rightarrow$  ~~given~~ efficiently computable  
(sq. roots modulo  $n = pq$ ).

• if  $\phi = 1 \bmod 4$ , no known poly-time deterministic  
algm to compute sq. roots of quadratic residues  
modulo  $\phi$ . (There is a poly-time Las Vegas algm,  
however)

(26)

(27)

Example:  $n = 77 = 7 \times 11$

$$e_k(x) = x^2 \pmod{77}$$

$$d_k(y) = \sqrt{y} \pmod{77}$$

To decrypt  $y=23$

$$7 = 4 \times 1 + 3, \quad 11 = 4 \times 2 + 3$$

$$\therefore (23)^{\frac{7+1}{4}} = (23)^{\frac{1}{4}} \pmod{7}$$

$$(23)^{\frac{7+1}{4}} = (23)^{\frac{1}{4}} \pmod{7} = 2^2 \pmod{7} = 4 \pmod{7}$$

$$(23)^{\frac{11+1}{4}} = (23)^{\frac{3}{4}} \pmod{11} = 2^3 \pmod{11} = 8 \pmod{11}$$

$$= 1^3 \pmod{11} = 1 \pmod{11}$$

Apply CRT

$$\begin{cases} x \equiv 4 \pmod{7} \\ x \equiv 1 \pmod{11} \end{cases}$$

if  $a^{-1} \pmod{b} = 1$   
 $a^{-1} u \equiv 1 \pmod{b}$   
 $\Rightarrow u \equiv 1 \pmod{b}$   
 $\Rightarrow u = a^{-1} \pmod{b}$

compute  
using  
Extended  
Euclidean  
algm.

$$11x^2 \equiv 1 \pmod{7}$$

$$11x^2 \equiv 1 \pmod{11}$$

$$7x^2 \equiv 1 \pmod{8}$$

$$\text{Soln. } \pm 10, \pm 32 \pmod{77}$$

Ideg  $x = 4 \times 11 \times y_1 + 1 \times 7 \times y_2 \pmod{77}$

$$\begin{aligned} &= 4 \times 11 \times 2 \\ &\quad + 1 \times 7 \times 8 \\ &= 67 \pmod{77} \\ &= -1 \pmod{77} \end{aligned}$$

$$\begin{aligned} &11 \pmod{7} \\ &\text{1 mod 7} \\ &7 \pmod{11} \\ &1 \pmod{8} \end{aligned}$$

$$\begin{aligned} &\text{gcd}(7, 11) = 1 \\ &\Rightarrow 7u + 11v = 1 \\ &\Rightarrow 7u \equiv 1 \pmod{11} \\ &\Rightarrow u \equiv 1 \pmod{11} \end{aligned}$$

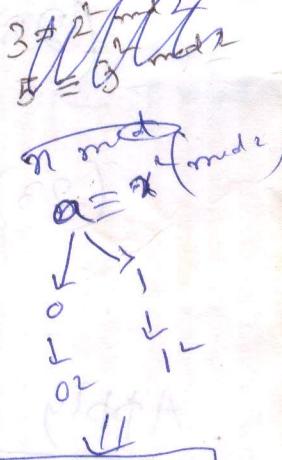
Disadvantage:

- encryption fun.  $E_k$  is not injective, so decryption cannot be done in an unambiguous fashion.
- for a valid ciphertext  $y$  ( $= x^2 \pmod n$  for plaintext  $x \in \mathbb{Z}_n^*$ ), there are 4 ~~four~~ sq. roots of  $y$  modulo  $n$ .  
 $\Rightarrow$  4 possible plaintexts that encrypt to  $y$ .
- No way for the decryptor to distinguish which of these 4 possible plaintexts is the 'right' plaintext. (unless the plaintext contains sufficient redundancy to eliminate 3 of these possible values).

Finding

$$\text{Sq. root mod } n = pqr \text{ when } p=2$$

$$a \equiv x^2 \pmod{2q} \text{ iff } \begin{cases} a \equiv x^2 \pmod{2} \\ \text{and } a \equiv x^2 \pmod{q} \end{cases}$$

Example:

Find  $\sqrt{68} \pmod{86}$ .  $\Downarrow$   
 $\sqrt{68} \pmod{86}$  is the sq. root of  $a$  modulo  $n=2q$ .

Ans:  $\sqrt{38} \pmod{43}$ .  $\Downarrow$   
 $\sqrt{38} \pmod{43}$  given  $a$ , find  $x$ .

All integers have unique sq. root mod 2.

• Check if  $a \in \{0^2, 1^2, 2^2, \dots, (\frac{q-1}{2})^2\} \pmod{q}$ .

- if  $a \equiv x^2 \pmod{q}$ , then  $a$  has unique sq. root mod  $2q$   
 $\Downarrow$  obtained by CRT to solve the system

$$\text{i.e. } \sqrt{a} \equiv 0 \pmod{q} \quad x \equiv 0 \pmod{q}$$

$$x \equiv \sqrt{a} \pmod{2}$$

- if  $a \equiv x^2 \pmod{q}$ ,  $x \neq 0$ , then  $a$  has two sq. roots mod  $2q$   
 $\Downarrow$  obtained by using CRT to solve the system

$$\begin{cases} x_1 \equiv 0 \pmod{q} \\ x_2 \equiv \sqrt{a} \pmod{q} \end{cases} \text{ and } \begin{cases} x_1 \equiv -\sqrt{a} \pmod{q} \\ x_2 \equiv \sqrt{a} \pmod{q} \end{cases}$$

$$n \equiv \sqrt{a} \pmod{2}$$

Example: find  $\sqrt{68} \pmod{86}$

6  
28

Soln.  $86 = 2 \times 43$

$$\left. \begin{array}{l} 68 = x^2 \pmod{2} \\ 68 = x^2 \pmod{43} \end{array} \right\}$$

$$68 = x^2 \pmod{2} \Rightarrow x = 0 \pmod{2}$$

$$68 = x^2 \pmod{43} \Rightarrow x = ? \pmod{43}$$

Check whether  $68 \in \left\{ 0^2, 1^2, 2^2, \dots, \left(\frac{43-1}{2}\right)^2 \right\}$ .  $\rightarrow$  expensive

Note that—  $43 = 4 \cdot 10 + 3$

$\therefore x = \pm 68^{\frac{43+1}{4}}$  are sq. root of  $68 \pmod{843}$ .

$$= \pm 68^{11}$$

Calculate  $68^{11}$  (use fast exponentiation)

$$68 = 25 \pmod{43}$$

$$(68)^2 = 25 \times 25 = 23 \pmod{43}$$

$$(68)^2 = 23 \times 23 = 13 \pmod{43}$$

$$(68)^3 = 13 \times 13 = 40 \pmod{43}$$

$$\therefore (68)^{11} = 25 \times 23 \times 40$$

$$= 38 \pmod{43}$$

$$x = \pm 38 \pmod{43}$$

Apply CRT

(20)

$$\therefore \begin{cases} x \equiv 0 \pmod{2} \\ x \equiv 38 \pmod{43} \end{cases}$$

$$\text{Solv} \rightarrow x = \left\{ 0 \times 43 \times (\overline{1} \pmod{2}) + 38 \times 2 \times (\overline{2} \pmod{43}) \right\} \pmod{86}$$

$$= (38 \times 2 \times 2) \pmod{86}$$

$$= 38 \pmod{86} \quad \left| \begin{array}{l} 22 \times 2 = 44 \\ = 1 \pmod{43} \\ \overline{2} \pmod{43} \\ = 22 \end{array} \right.$$

$$\therefore \begin{cases} x \equiv 0 \pmod{2} \\ x \equiv -38 \pmod{43} = \overline{5} \pmod{43} \end{cases}$$

$$\text{Solv} \rightarrow x = (\overline{5} \times 2 \times 22) \pmod{86}$$

$$= 48 \pmod{86}.$$

- fast modular exponentiation (to find high power).
- Extended Euclidean Algm. (to find modular inverse)

Example: find  $\sqrt{23} \pmod{77}$

To find  $x$

$$23 \equiv x^2 \pmod{7}$$

$$7 = 4 \cdot 1 + 3$$

$$23 \equiv x^2 \pmod{11}$$

$$11 = 4 \cdot 2 + 3$$

Sq. roots of  $23 \pmod{7}$   $\rightarrow$

$$x \equiv \pm (23)^{\frac{7+1}{4}} \pmod{7}$$

Sq. roots of  $23 \pmod{11}$   $\rightarrow$

$$x \equiv \pm (23)^{\frac{11+1}{4}} \pmod{7}$$

$$\text{i.e. } x \equiv \pm 4 \pmod{7}$$

$$x \equiv \pm 1 \pmod{11}$$

Apply CRT

$$\begin{cases} x \equiv 1 \pmod{7} \\ x \equiv 1 \pmod{11} \end{cases}$$

$$x = 4 \cdot 11 (\overline{1} \pmod{7}) + 1 \cdot 7 (\overline{7} \pmod{11})$$

$$= 4 \cdot 11 \cdot 2 + 1 \cdot 7 \cdot 8 = 67 \pmod{77} = -10 \pmod{77}$$

$\sqrt{y} \pmod{n}$

$$n = p q$$

$$p, q = 3 \pmod{4}$$

$$\left( \pm y^{\frac{p+1}{4}} \right)^2 = y^{\frac{p+1}{2}} = y^{\frac{p-1}{2}} \cdot y$$

$$= y \pmod{p}$$

as by Euler's criterion

$$\text{Legendre symbol} \left( \frac{y}{p} \right) = y^{\frac{p-1}{2}} \pmod{p}$$

$$= 1 \text{ if } y \in Q_p,$$

$\pm y^{\frac{p+1}{2}}$  after two sq. roots  
of  $y \pmod{p}$ .

## Probabilistic encryption

(30/1)

- encryption is made probabilistic.
  - many possible encryptions of each plaintext
  - not feasible to test whether a given ciphertext is an encryption of a particular plaintext
  - no information about plaintext should be computable from the ciphertext (in poly. time).

Decisional Composite residuosity assumption

## Goldwasser-Micalli Probabilistic Public-key encryption

(AFL bit encryption)

- $n = p \cdot q$ ,  $p, q$  are primes
- $m \in \widetilde{QR}(n)$  i.e.  $\left(\frac{m}{p}\right) = \left(\frac{m}{q}\right) = -1$ . secret  
not a square mod n
- $P = \{0, 1\}$ ,  $C = \mathbb{Z}_n^*$ ,  $K = \{(n, p, q, m) : n = p \cdot q$ ,  
 $p, q$  primes,  $m \in \widetilde{QR}(n)\}$
- $p, q \rightarrow \text{secret}$ ,  $m \rightarrow \text{public}$

Encryption      choose  $y \in \mathbb{Z}_n^*$ , message  $x \in \{0, 1\}$

$$PK = (n, m)$$

$$SK = (p, q)$$

$$e_{PK}(x, y) = m^x \cdot y^2 \pmod{n}$$

Decryption       $d_{SK}(y) = \begin{cases} 0 & \text{if } y \in QR(n) \\ 1 & \text{if } y \notin \widetilde{QR}(n) \end{cases}$

Correctness      On receiving  $y \in QR(n) \cup \widetilde{QR}(n)$ , decryptor can use his knowledge of factorization of  $n$  to determine whether  $y \in QR(n)$  or whether  $y \in \widetilde{QR}(n)$  by computing:

$$\left(\frac{y}{p}\right) = y^{\frac{p-1}{2}} \pmod{p}$$

(31) (8)

$$\text{then } y \in QR(n) \Leftrightarrow \left(\frac{y}{p}\right) = 1.$$

$$n=pq \quad J_n = \left\{ a \in \mathbb{Z}_n^* \mid \left(\frac{a}{n}\right) = 1 \right\}$$

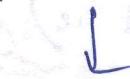
$$= QR(n) \cup \tilde{QR}(n)$$

squares mod n



$$\left(\frac{a}{p}\right) = \left(\frac{a}{q}\right) = 1$$

pseudosquare mod n.



$$\left(\frac{a}{p}\right) = \left(\frac{a}{q}\right) = -1.$$

$$y \in QR(n) \cup \tilde{QR}(n) \Rightarrow \left(\frac{y}{n}\right) = 1.$$

$$\therefore y \in QR(n) \Leftrightarrow \left(\frac{y}{p}\right) = 1.$$

Security: Given  $n=pq$  &  $\left(\frac{y}{n}\right) = 1$

(Dec. Composite Residuosity)   
 Decide whether  $y \in QR(n)$  or not

- Homo morphic property if  $C_0, G$   $\rightarrow$  encryption of  $m_0, m_1$
- Applicable to design more complex crypto system. Then  $C_0, G$   $\rightarrow$  encryption of  $m_0 \oplus m_1$