# **Graph Theory: Euler Graphs and Digraphs**



#### Pallab Dasgupta,

Professor, Dept. of Computer Sc. and Engineering, IIT Kharagpur pallab@cse.iitkgp.ernet.in

Indian Institute of Technology Kharagpur

#### PALLAB DASGUPTA

### **Euler Circuit**

- We use the term *circuit* as another name for *closed trail*.
  - A circuit containing every edge of G is an *Eulerian circuit*.
  - A graph whose edges comprise a single closed trail is *Eulerian*.

### **Properties**

• Non-trivial maximal trails in even graphs are closed.

• A finite graph G is Eulerian if and only if all its vertex degrees are even and all its edges belong to a single component.

 For a connected nontrivial graph with 2k odd vertices, the minimum number of pairwise edge-disjoint trails covering the edges is max{k, 1}.

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Input: A graph G with one non-trivial component and at most two odd vertices.

Initialization: Start at a vertex that has odd degree unless G is even, in which case start at any vertex.

Iteration:

From the current vertex, traverse any remaining edge whose deletion from the graph does not leave a graph with two nontrivial components. Stop when all edges have been traversed.

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#### **Euler Trails in Directed Graphs**

- Input: A digraph G that is an orientation of a connected graph and has  $d^+(u) = d^-(u)$  for all  $u \in V(G)$ .
- **Step1:** Choose a vertex  $v \in V(G)$ . Let G' be the digraph obtained from G by reversing direction on each edge. Search G' to construct T' consisting of paths from v to all other vertices.
- **Step2:** Let T be the reversal of T'. T contains a *u*,*v*-path in G for each  $u \in V(G)$ . Specify an arbitrary ordering of the edges that leave each vertex *u*, except that for  $u \neq v$ , the edge leaving *u* in T must come last.
- Step3: Construct an Eulerian circuit from v as follows. Whenever u is the current vertex, exit along the next unused edge in the ordering specified for edges leaving u.

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## **The Chinese Postman Problem**

- Suppose a mail carrier traverses all edges in a road network, starting and ending at the same vertex.
  - The edges have non-negative weights representing distance or time.
  - We seek a closed walk of minimum total length that uses all the edges.

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