

Problem Reduction Search: AND/OR Graphs & Game Trees

Course: CS40002

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Problem Reduction Search

- Planning how best to solve a problem that can be recursively decomposed into sub-problems in multiple ways
 - ◆ Matrix multiplication problem
 - ◆ Tower of Hanoi
 - ◆ Blocks World problems
 - ◆ Theorem proving

Formulations

■ AND/OR Graphs

- ◆ An OR node represents a choice between possible decompositions
- ◆ An AND node represents a given decomposition

■ Game Trees

- ◆ Max nodes represent the choice of my opponent
- ◆ Min nodes represent my choice

The AND/OR graph search problem

- **Problem definition:**
 - **Given:** $[G, s, T]$ where
 - G : implicitly specified AND/OR graph
 - S : start node of the AND/OR graph
 - T : set of terminal nodes
 - $h(n)$ heuristic function estimating the cost of solving the sub-problem at n
 - **To find:**
 - A minimum cost solution tree

Algorithm AO*

- 1. Initialize:** **Set $G^* = \{s\}$, $f(s) = h(s)$**
 If $s \in T$, label s as SOLVED
- 2. Terminate:** **If s is SOLVED, then Terminate**
- 3. Select:** **Select a non-terminal leaf node n**
 from the marked sub-tree
- 4. Expand:** **Make explicit the successors of n**
 For each new successor, m :
 Set $f(m) = h(m)$
 If m is terminal, label m SOLVED
- 5. Cost Revision:** **Call $\text{cost-revise}(n)$**
- 6. Loop:** **Go To Step 2.**

Cost Revision in AO*: cost-revise(n)

1. Create $Z = \{n\}$
2. If $Z = \{ \}$ return
3. Select a node m from Z such that m has no descendants in Z
4. If m is an AND node with successors

r_1, r_2, \dots, r_k :

Set $f(m) = \sum [f(r_i) + c(m, r_i)]$

Mark the edge to each successor of m

If each successor is labeled SOLVED,
then label m as SOLVED

Cost Revision in AO*: cost-revise(n)

5. If m is an OR node with successors

r_1, r_2, \dots, r_k :

Set $f(m) = \min \{ f(r_i) + c(m, r_i) \}$

Mark the edge to the best successor of m

If the marked successor is labeled

SOLVED, label m as SOLVED

6. If the cost or label of m has changed, then insert those parents of m into Z for which m is a marked successor

7. Go to Step 2.

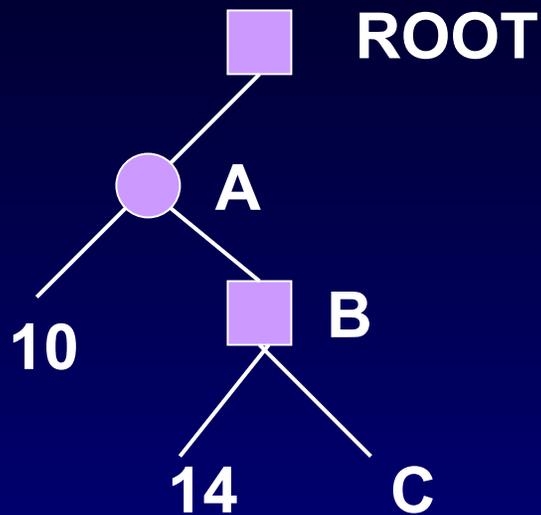
Searching OR Graphs

- How does AO* fare when the graph has only OR nodes?

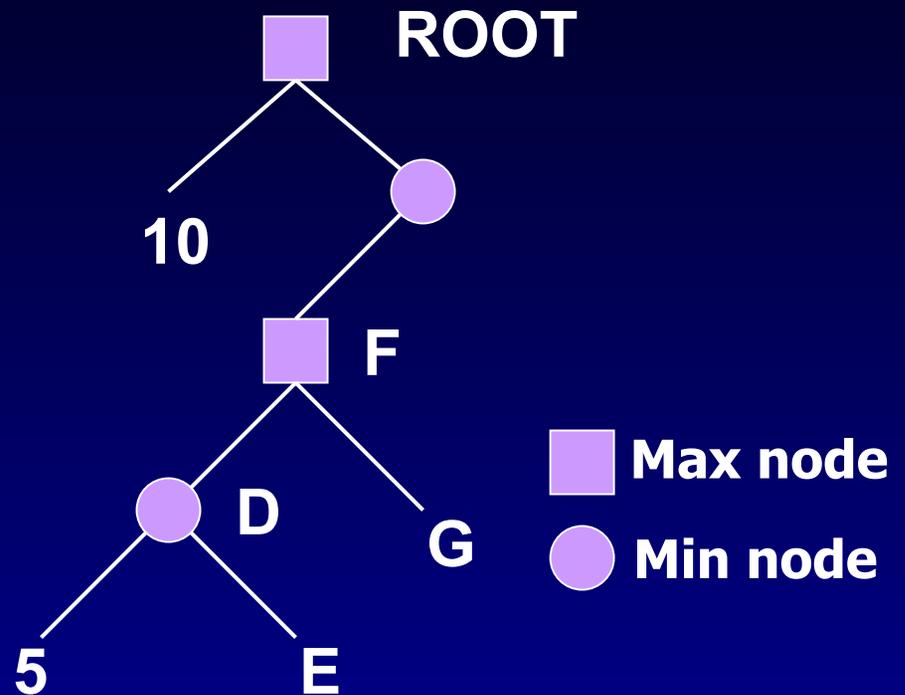
Searching Game Trees

- Consider an OR tree with two types of OR nodes, namely Min nodes and Max nodes
- In Min nodes, select the min cost successor
- In Max nodes, select the max cost successor
- **Terminal nodes are winning or losing states**
 - ◆ It is often infeasible to search up to the terminal nodes
 - ◆ We use heuristic costs to compare non-terminal nodes

Shallow and Deep Pruning



Shallow Cut-off



Deep Cut-off

Alpha-Beta Pruning

- Alpha Bound of J:
 - ◆ The max current val of all MAX ancestors of J
 - ◆ Exploration of a min node, J, is stopped when its value equals or falls below alpha.
 - ◆ In a min node, we update beta
- Beta Bound of J:
 - ◆ The min current val of all MIN ancestors of J
 - ◆ Exploration of a max node, J, is stopped when its value equals or exceeds beta
 - ◆ In a max node, we update alpha
- In both min and max nodes, we return when $\alpha \geq \beta$

Alpha-Beta Procedure: $V(J; \alpha, \beta)$

1. If J is a terminal, return $V(J) = h(J)$.
2. If J is a max node:
 - For each successor J_k of J in succession:
 - Set $\alpha = \max \{ \alpha, V(J_k; \alpha, \beta) \}$
 - If $\alpha \geq \beta$ then return β , else continue
 - Return α
3. If J is a min node:
 - For each successor J_k of J in succession:
 - Set $\beta = \min \{ \beta, V(J_k; \alpha, \beta) \}$
 - If $\alpha \geq \beta$ then return α , else continue
 - Return β