

Process

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and Mainack Mondal (even section)
CS39002

Spring 2019-20



Today's class

- What is a process?
- Structure of a process
- Process states
- Process control block
- Context switch

What is a process?

- A process is a program in execution
 - Recall multitasking
 - Several processes may be in various stages of execution at same time

What is a process?

- A process is a program in execution
 - Recall multitasking
 - Several processes may be in various stages of execution at same time
- CPU switches between several processes
 - We say that processes are executing concurrently
 - CPU multiplexed

Job scheduling and process scheduling

- Job scheduling
 - In batch processing or multiprogramming user programs are called jobs
 - Long term scheduler
- Process scheduling
 - Short term scheduler

Job scheduling and process scheduling

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Check the board

Today's class

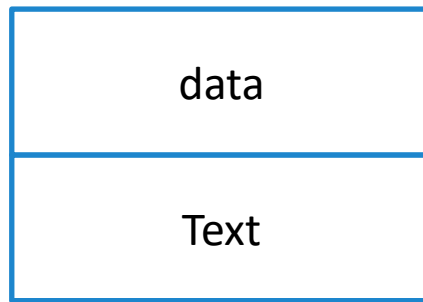
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What are parts of a process?



Program code

What are parts of a process?



Global variable

Program code

What are parts of a process?

heap
data
Text

Dynamically allocated memory

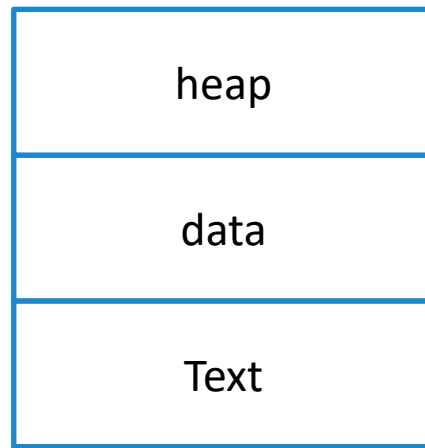
Global variable

Program code

What are parts of a process?



Temporary data,
function parameters,
local variables, return addresses

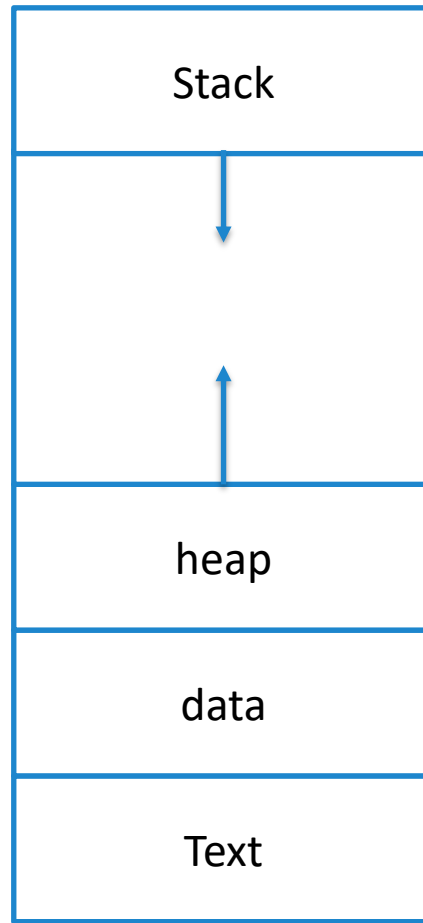


Dynamically allocated memory

Global variable

Program code

What are parts of a process?



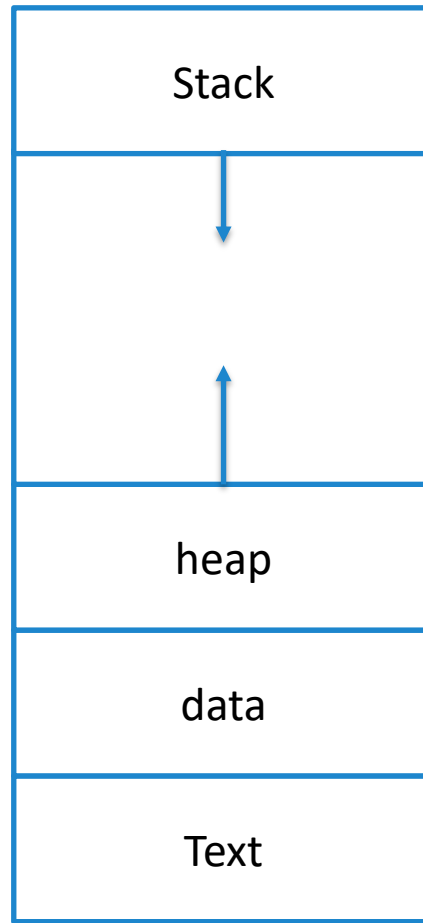
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What are parts of a process?



Temporary data,
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Also
Program
counter (PC),
CPU registers,
open files

Characteristics of a process

- Program is a passive entity
- Process is an active entity
 - Program becomes process when the code is loaded in the memory and ready to execute
- Each execution instance of the same program is a separate process

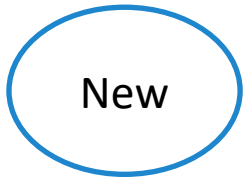
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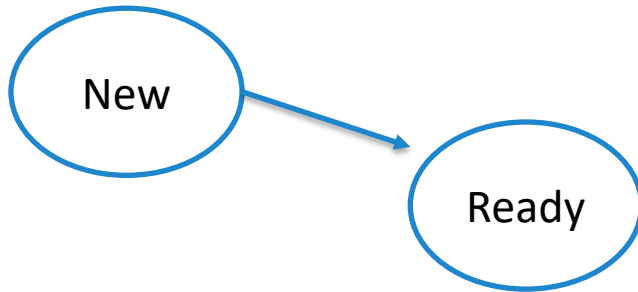
You can think process execution as an automata

- States
 - **new:** The process is being created
 - **running:** Instructions are being executed
 - **waiting:** The process is waiting for some event to occur
 - **ready:** The process is waiting to be assigned to a processor
 - **terminated:** The process has finished execution
- Additionally there is “swap space”
 - Resides in the disk
 - Swap-out and swap-in between main memory and disk

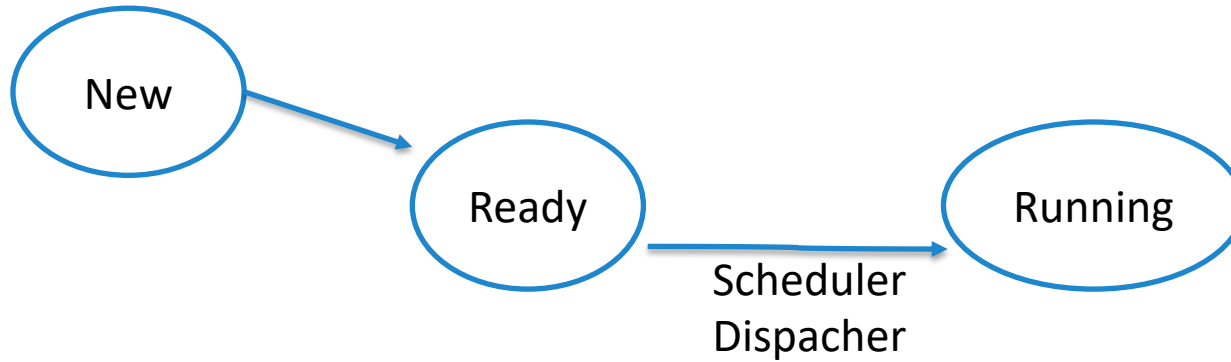
Process state diagram



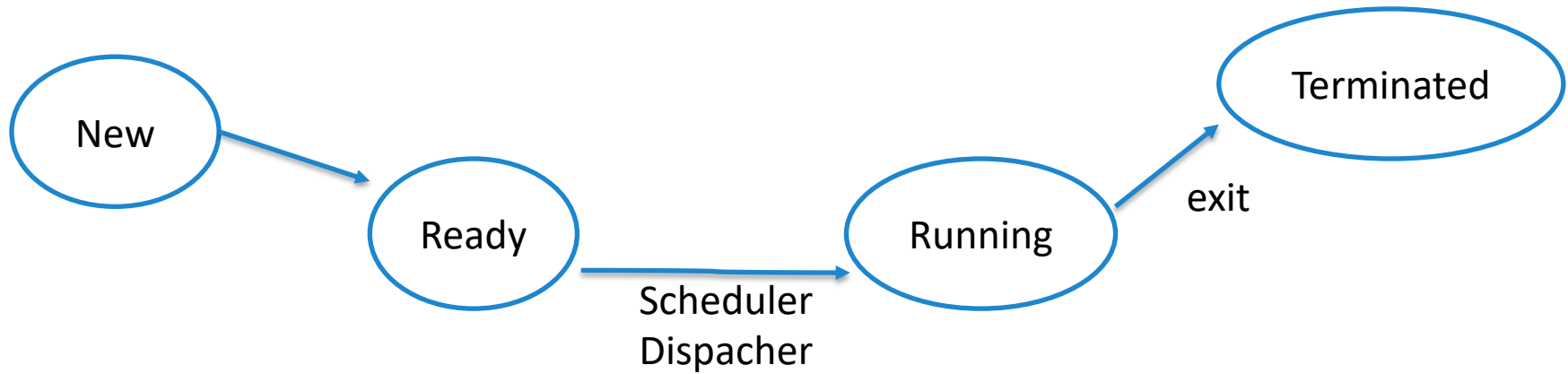
Process state diagram



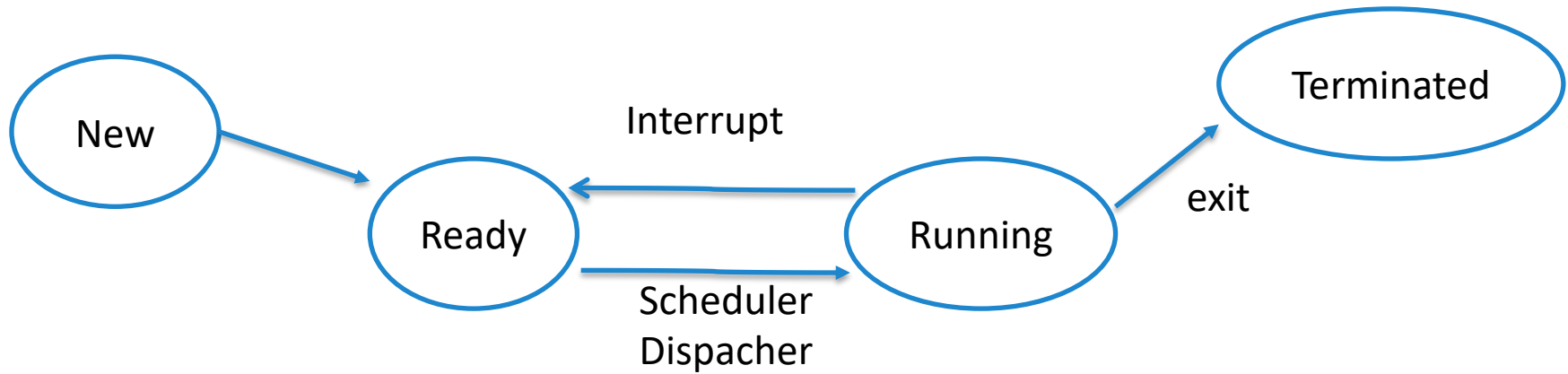
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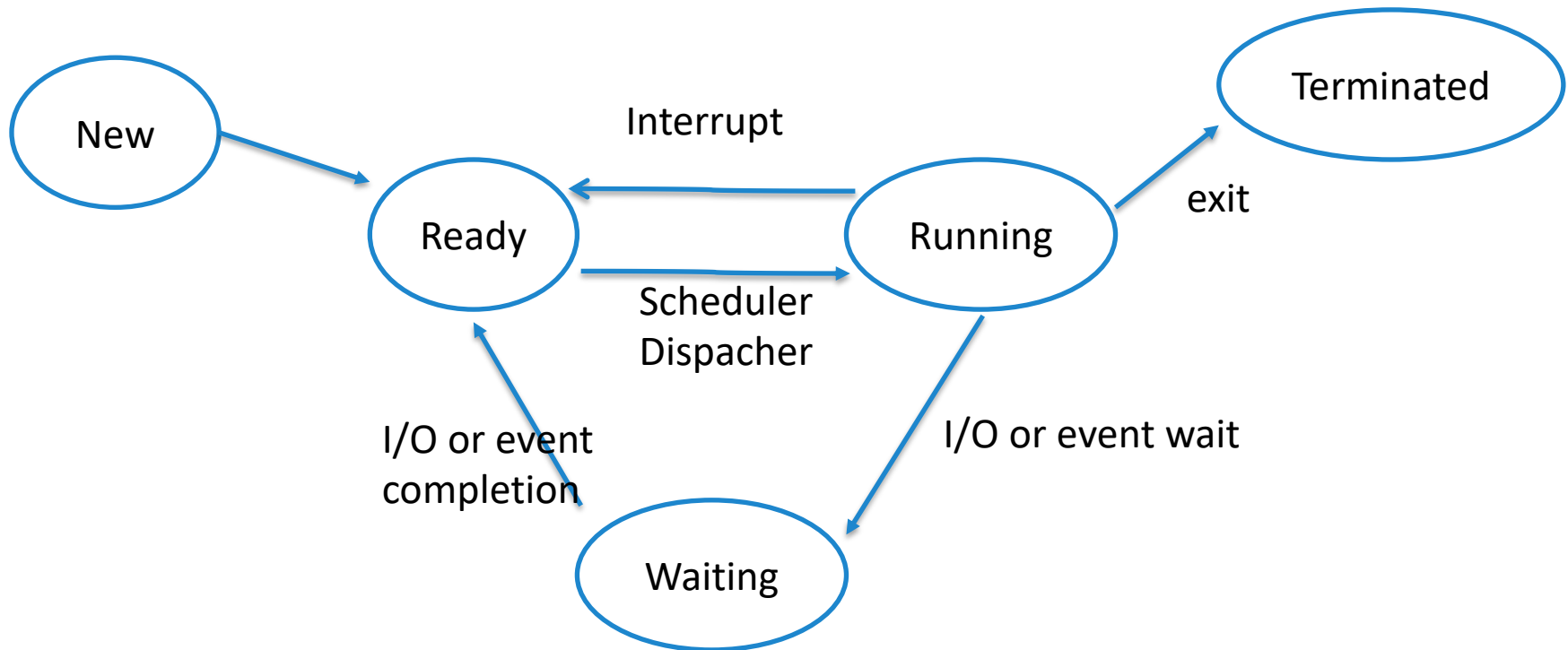
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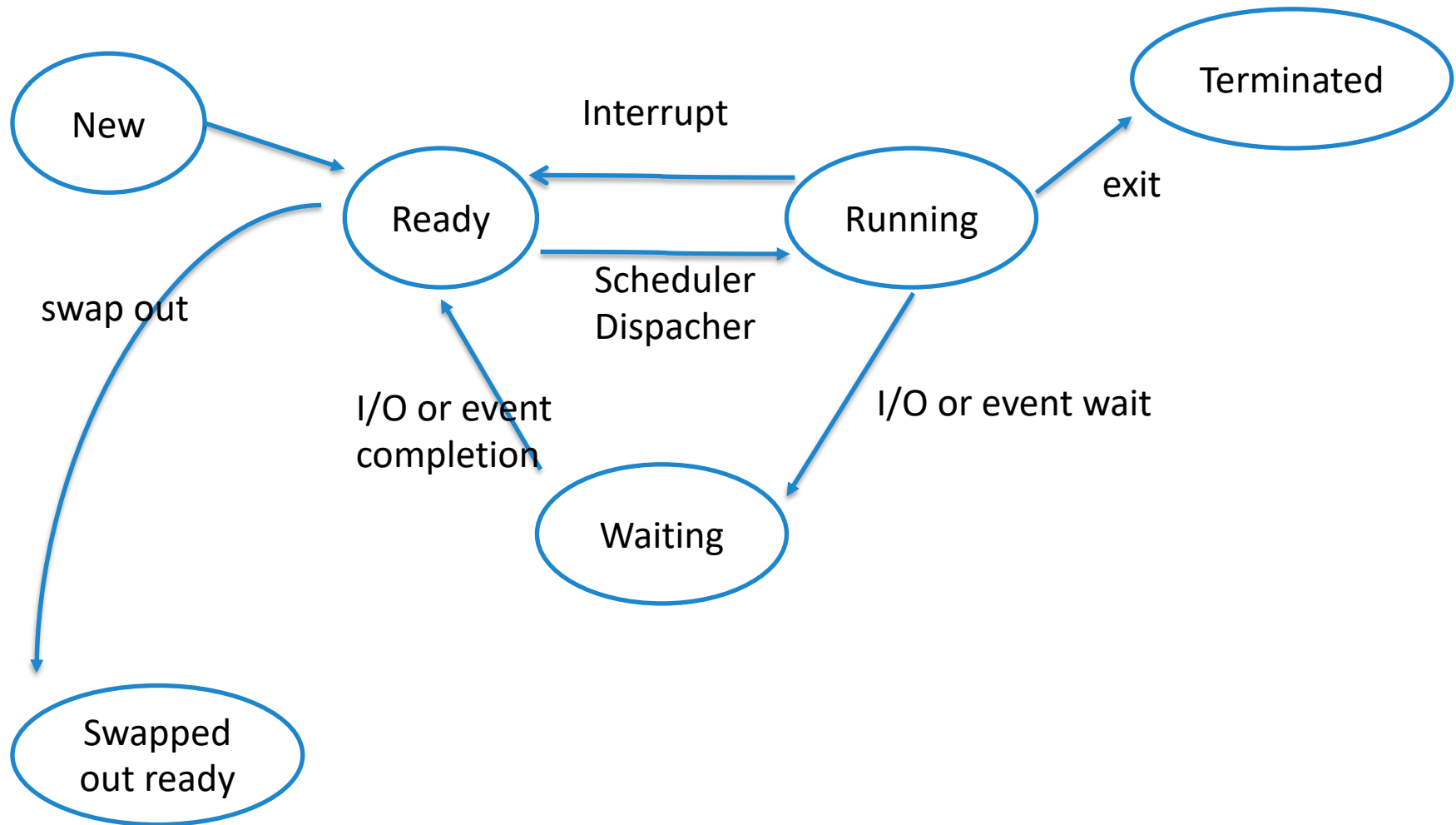
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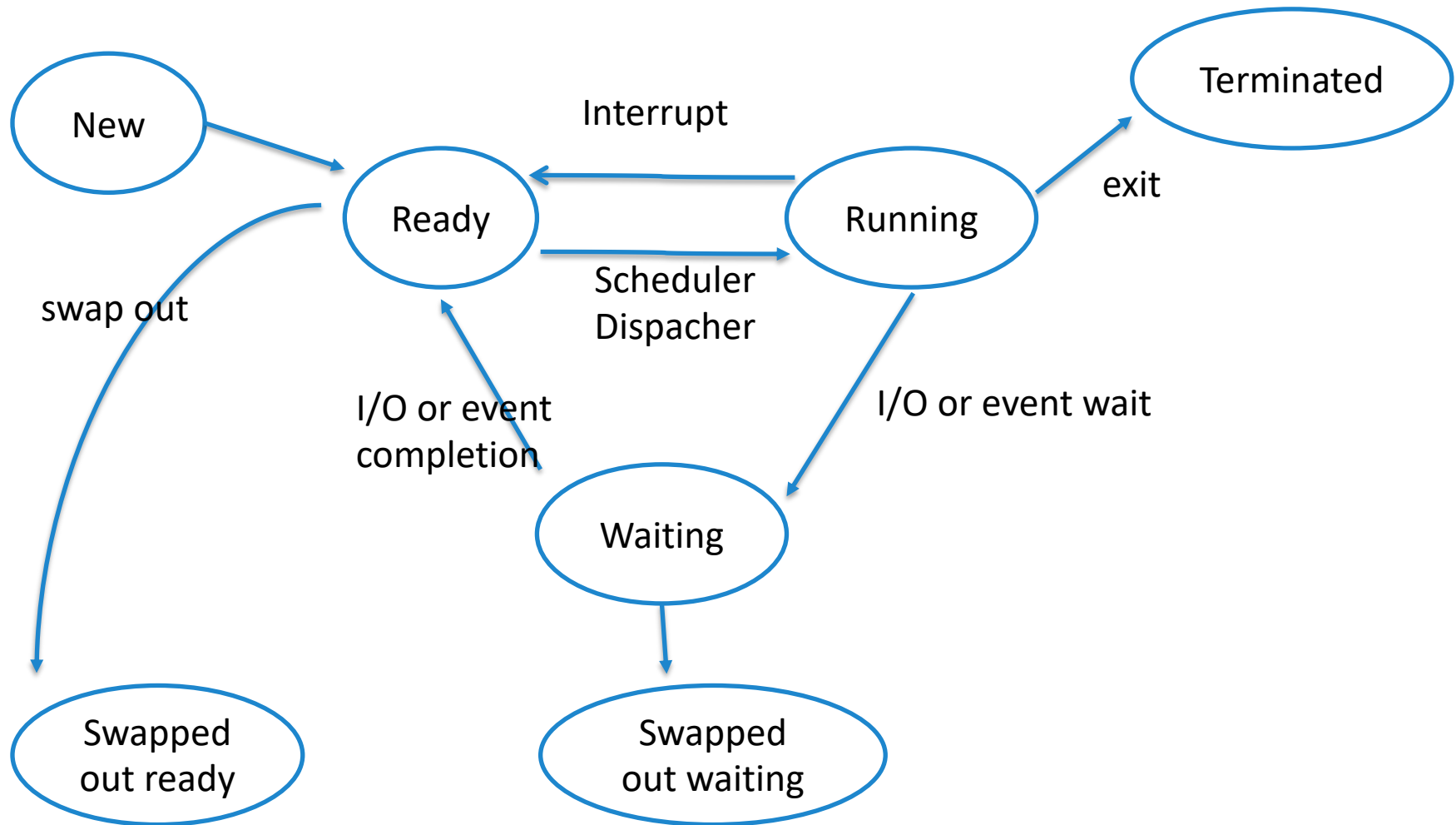
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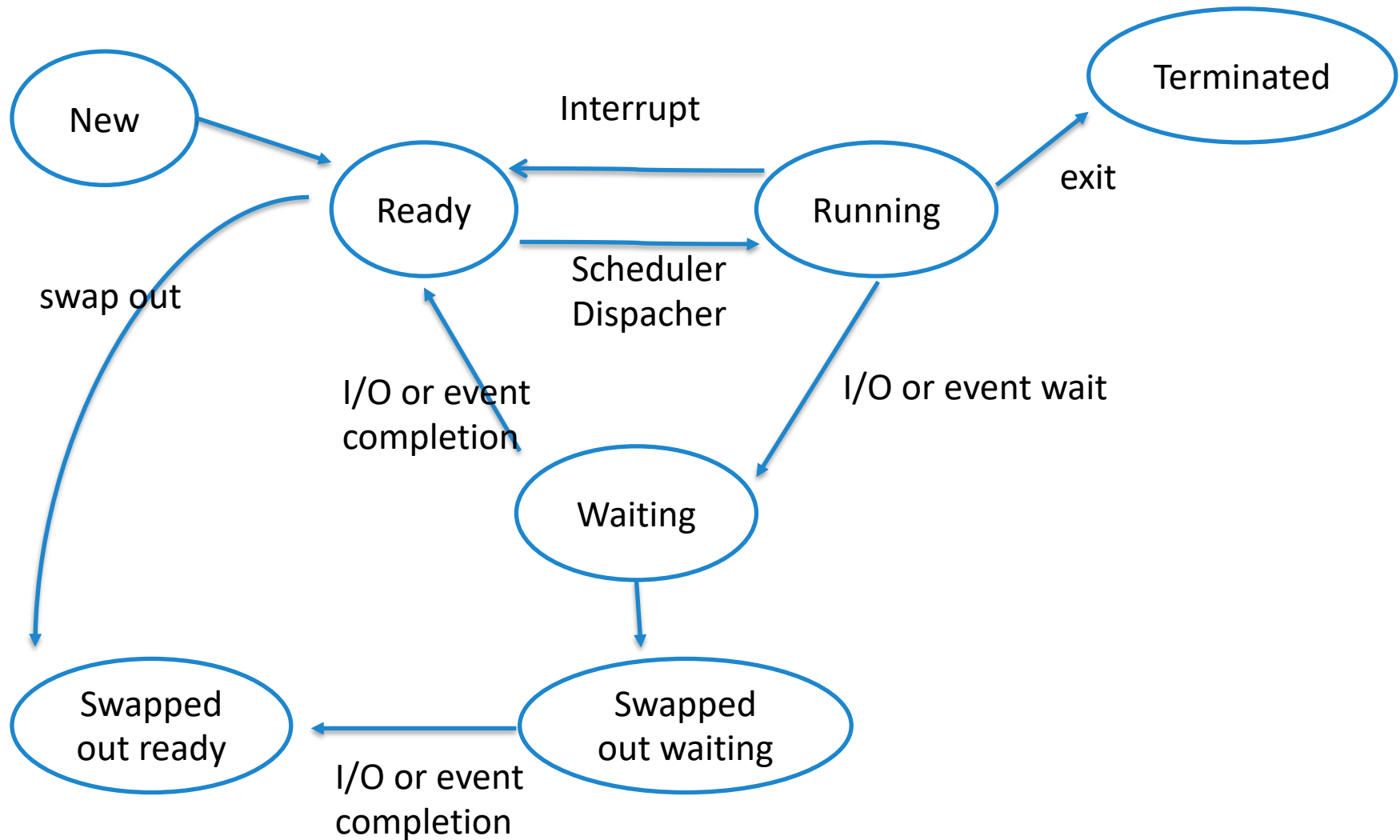
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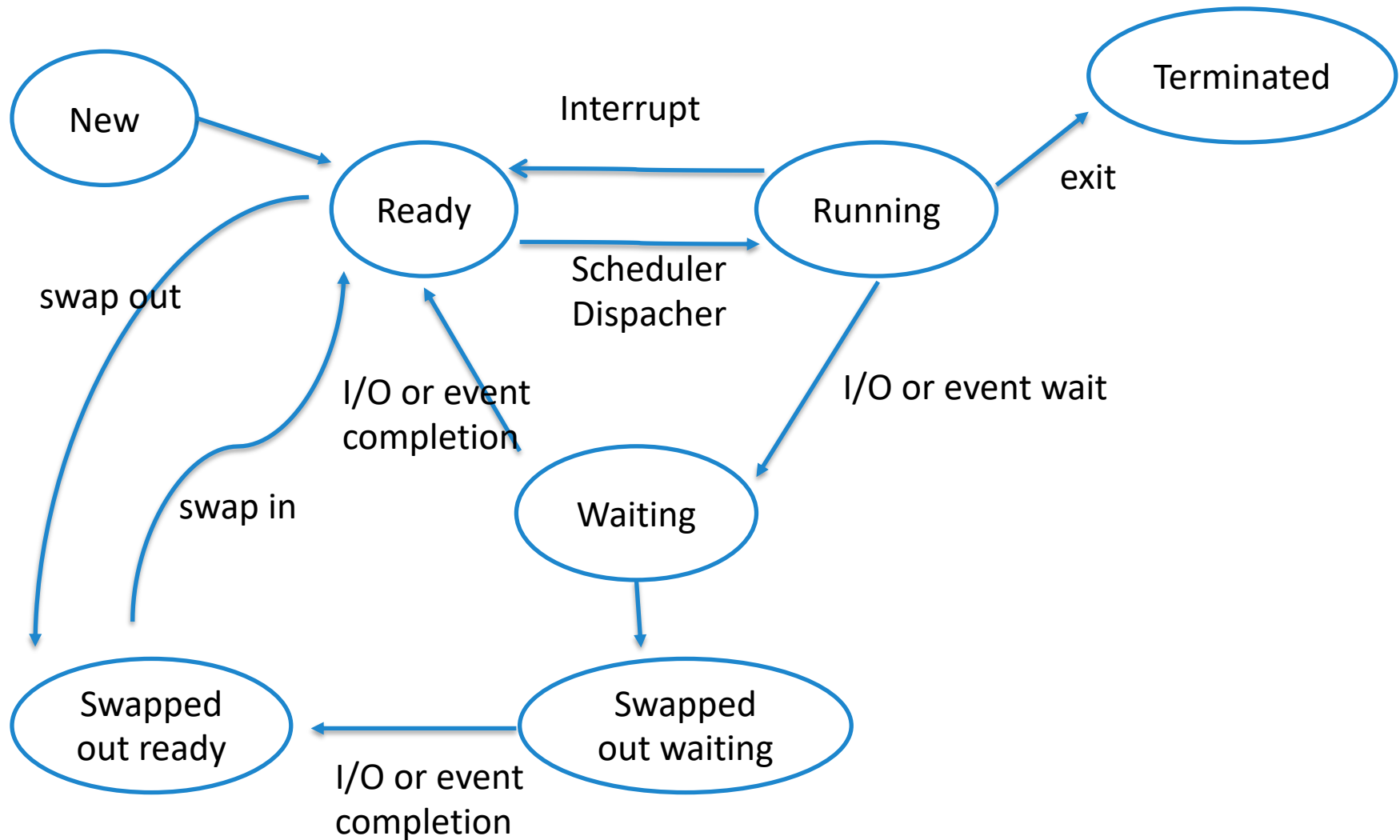
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Process control block (PCB)

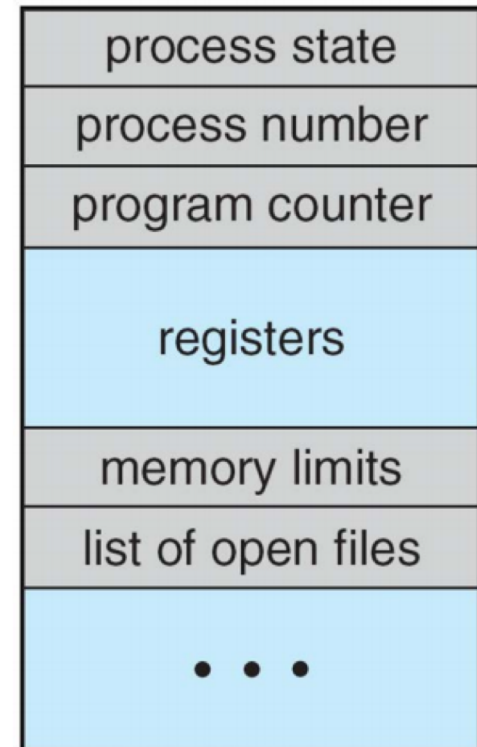
- Each process is represented in the kernel as a PCB
 - Also called task control block
 - Contains many pieces of information associated with a specific process

Structure of a PCB

- **Process state** – running, waiting, etc
- **Program counter (PC)** – location of instruction to next execute
- Content of CPU registers
- **CPU scheduling information**- priorities, scheduling queue pointers
- **Memory-management information** –Base and limit registers, page tables
- **Accounting information** – CPU used, clock time elapsed since start, time limits, pid
- **I/O status information** – I/O devices, allocated to process, list of open files

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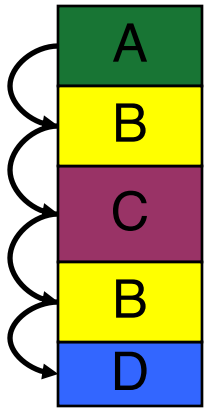


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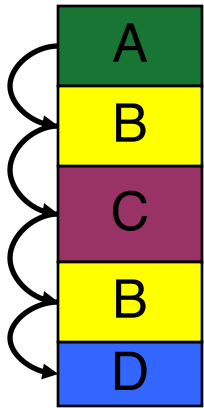
CPU runs multiple processes

CPU's point of view

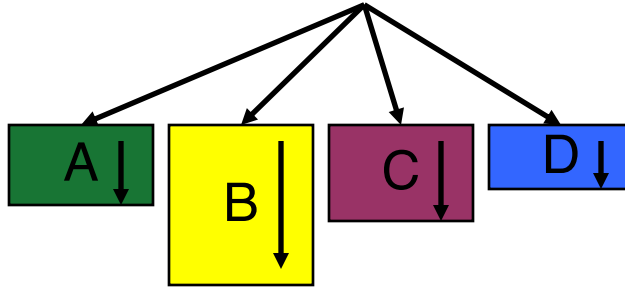


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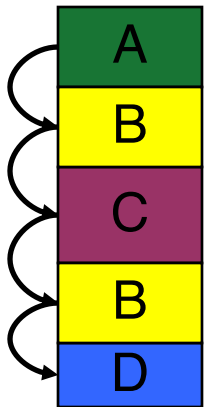


process point of view

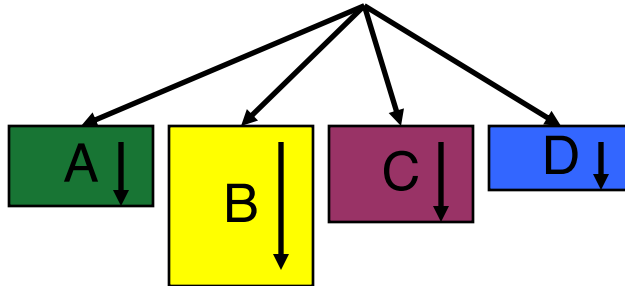


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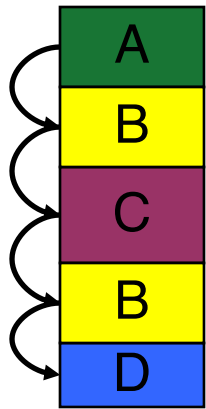
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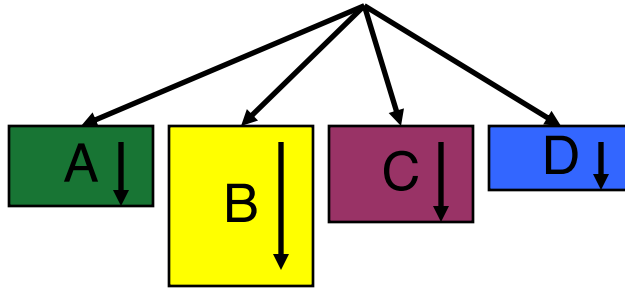
- Multiprogramming of four programs
- Conceptual model
 - 4 independent processes
 - Processes run sequentially
- Only one program active at any instant!
 - That instant can be very short...

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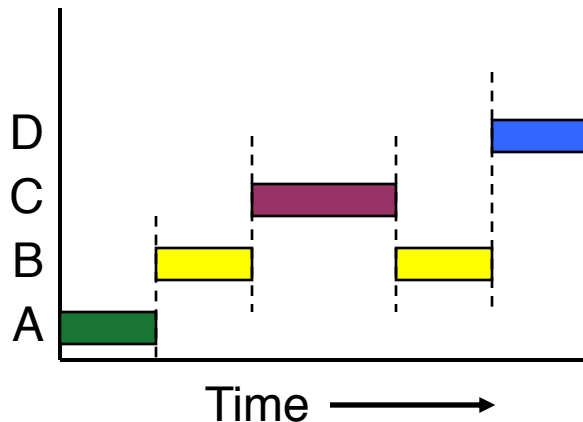
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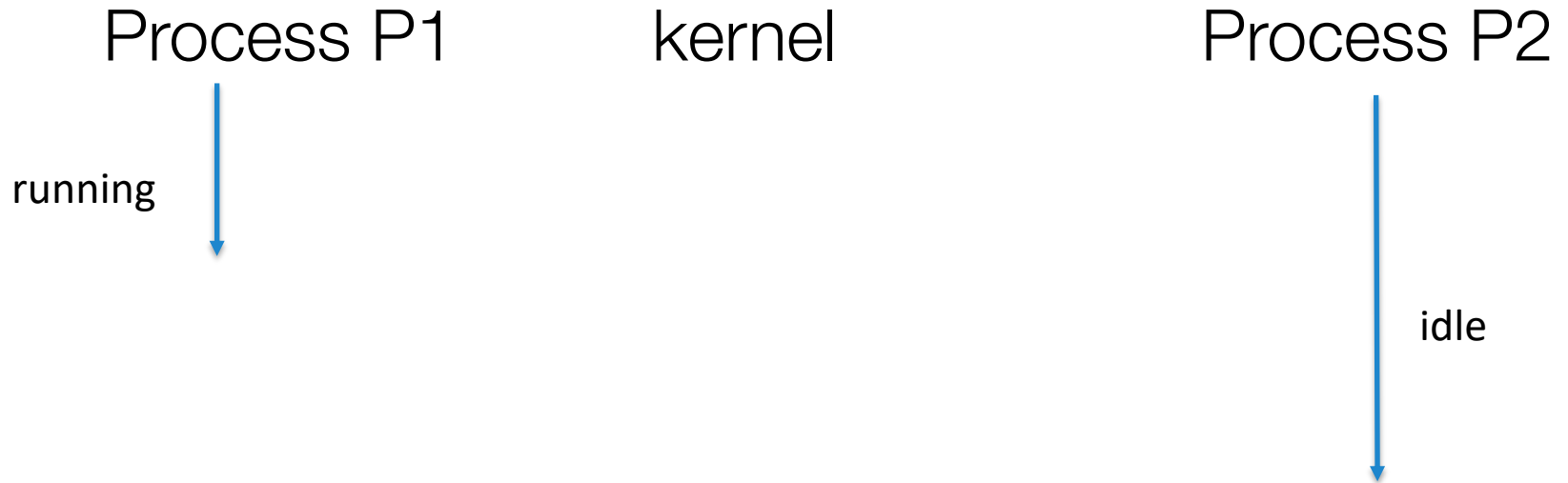


Gantt chart for multiprogramming

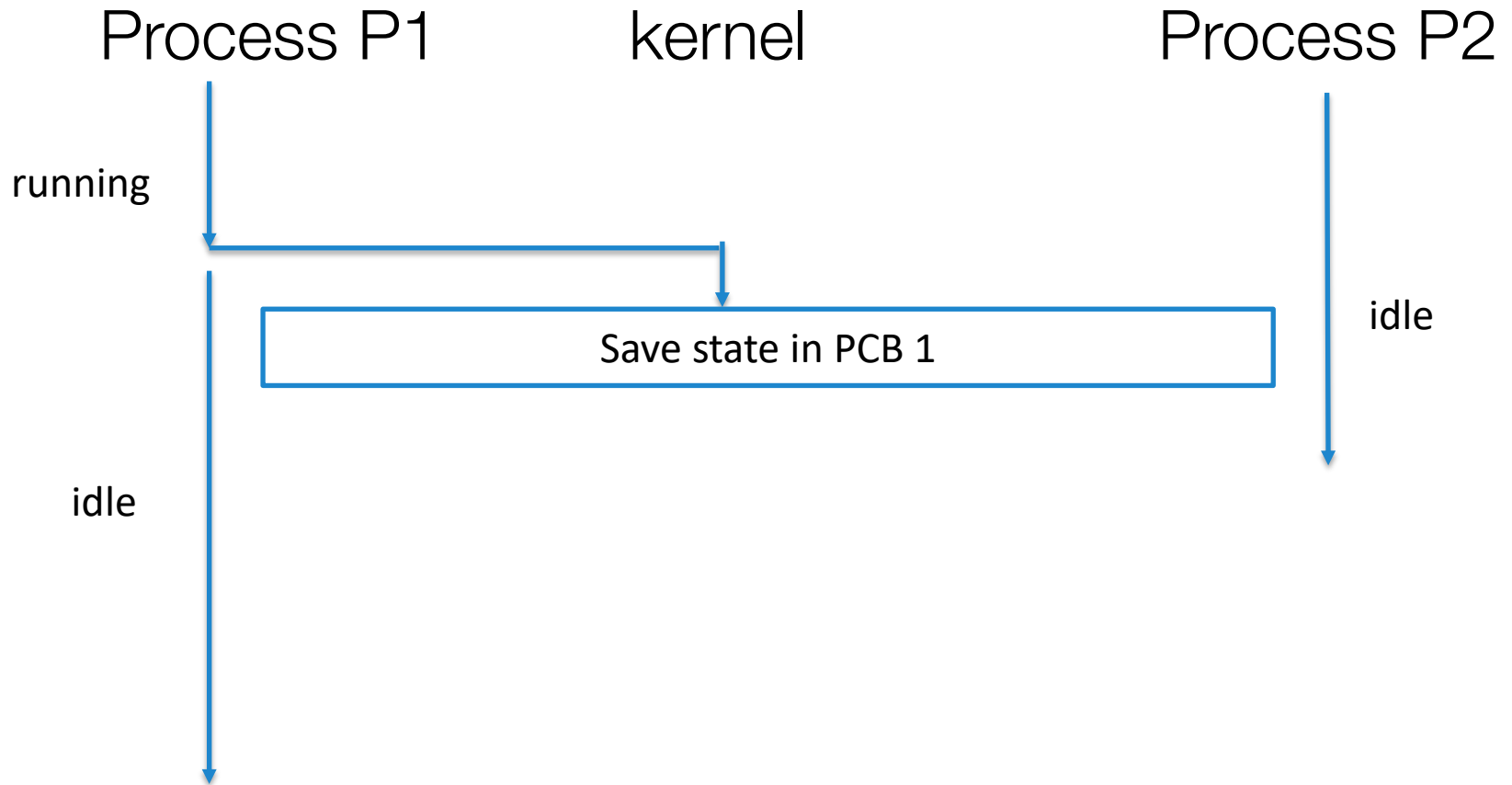
How to interleave processes?

- CPU switches to another process
 - the system saves the state of the old process and load the saved state for the new process via a context switch
 - Context of a process == PCB
 - More complex the OS and PCB, longer to switch

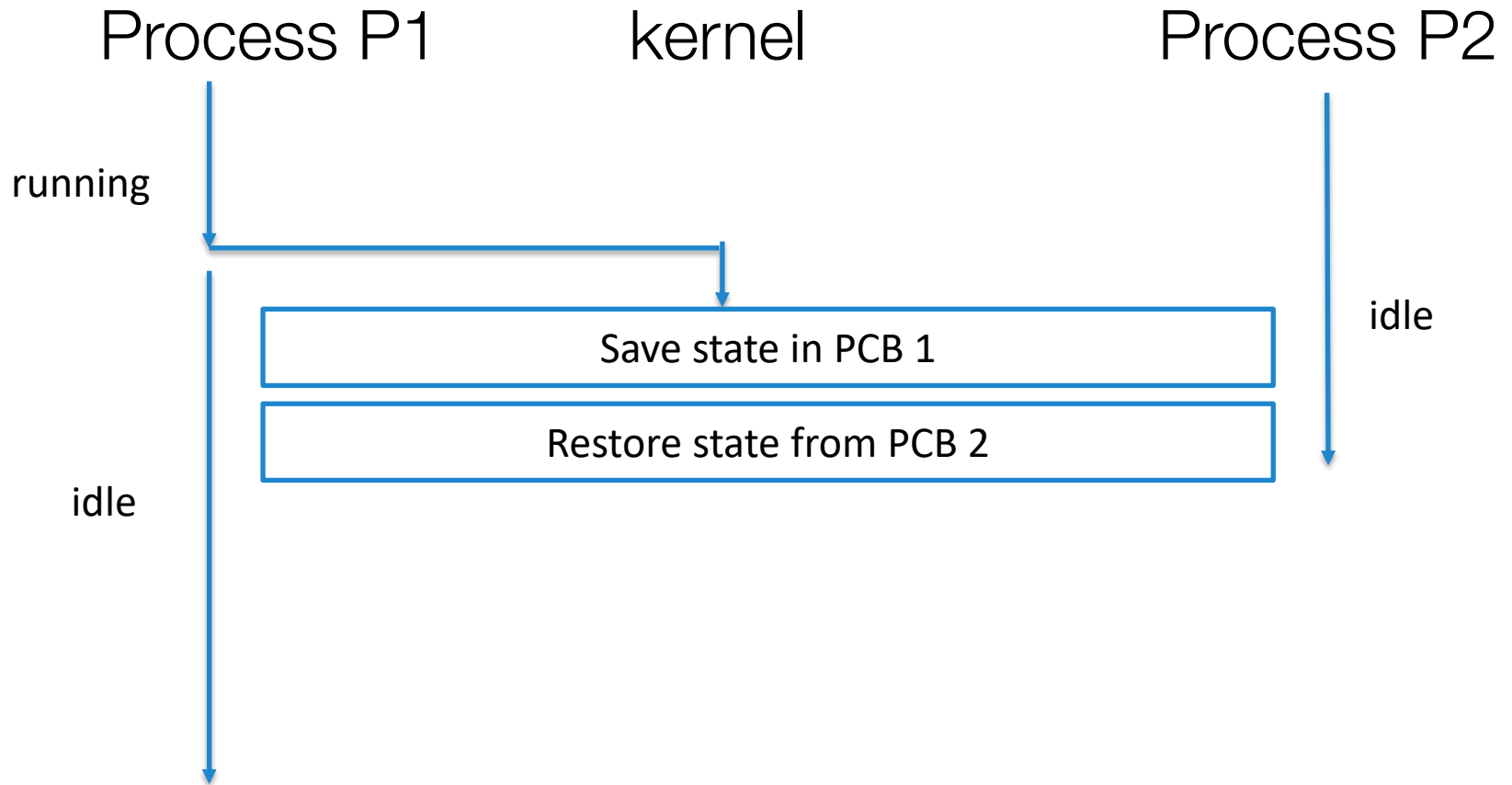
Context switch



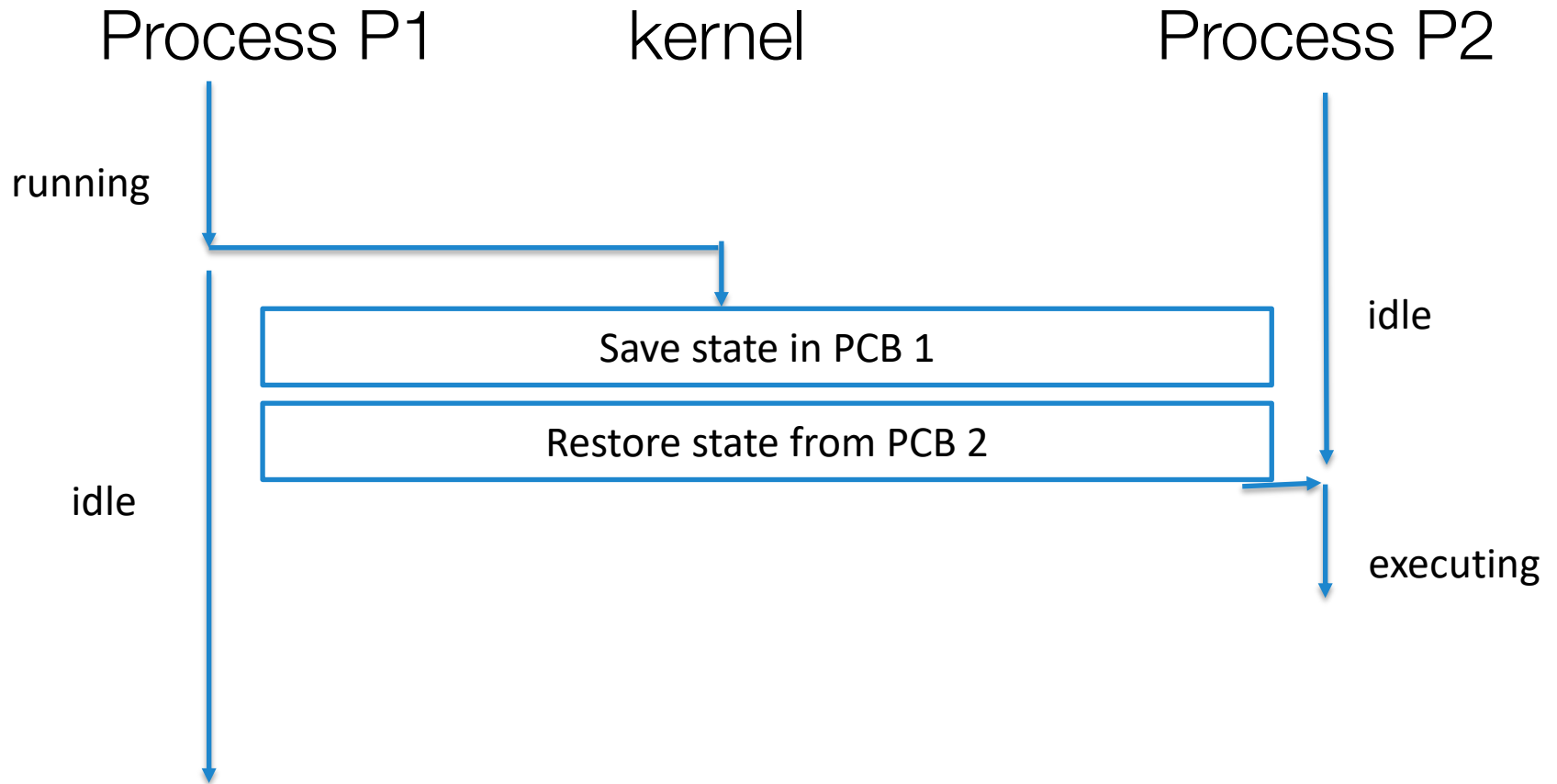
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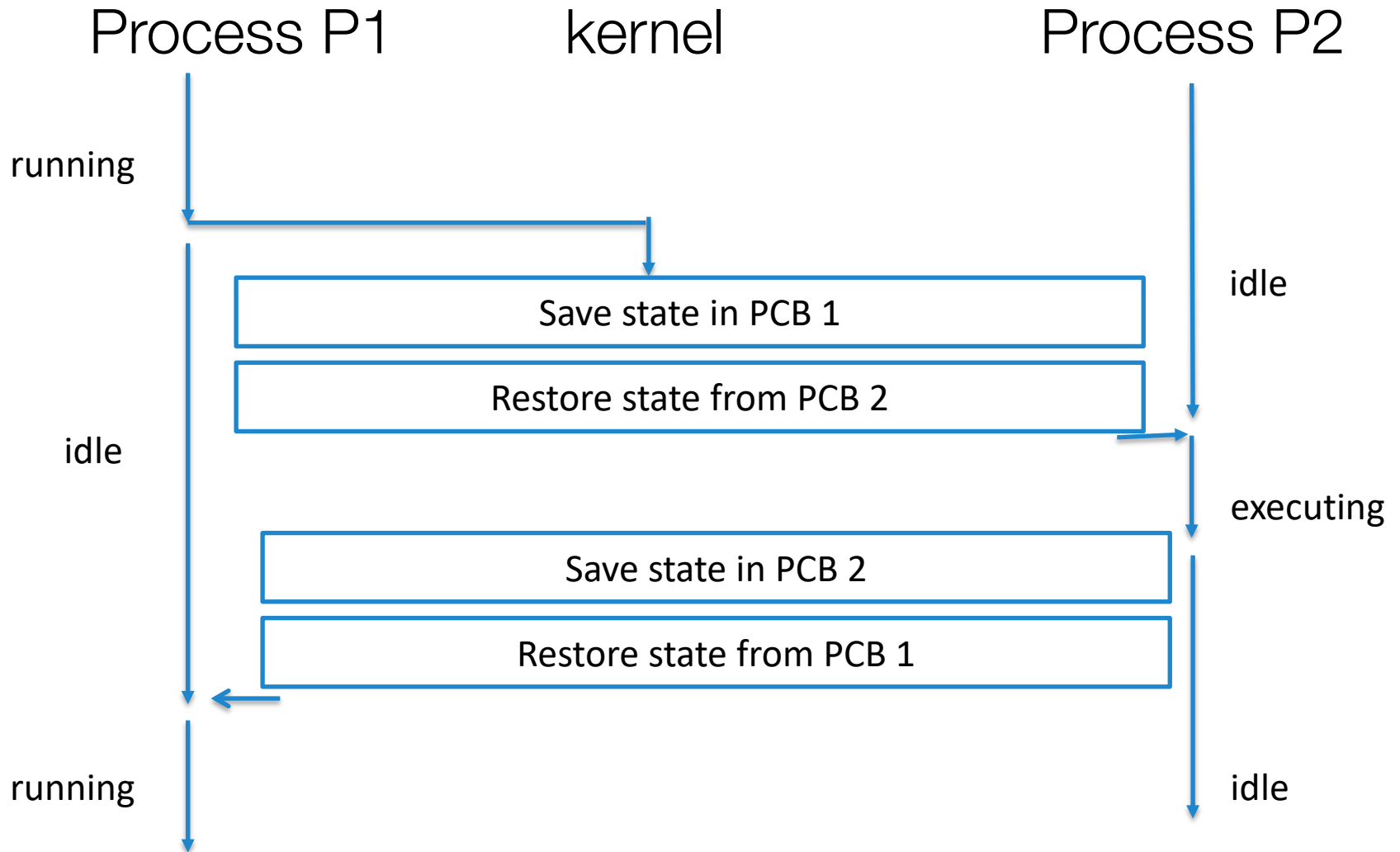
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Process scheduling and operations

- We will look at process scheduling the next day
- Today let's look at an operation

pipes()

- Acts as a medium to allow two processes to communicate
 - Communication can be uni/bi directional
 - Must there exist a relationship (i.e., parent-child) between the communicating processes?
 - Can the pipes be used over a network?

Ordinary pipes

- A message passing medium between related processes
 - Cannot be accessed from outside the process
 - Typically, a parent process creates a pipe and uses it to communicate with its child process
 - Pipes behave like FIFO queues
 - Read-write in pipe == producer-consumer
 - Producer writes to one end (the write-end of pipe)
 - Consumer reads from the other end (the read-end of pipe)
 - Unidirectional

Producer consumer in pipes

Producer

```
message next_produced;  
while(TRUE){  
    <produce and put data in next produced>  
    ...  
    send( next_produced);  
}
```

Consumer

```
message next_consumed;  
while(TRUE){  
    receive(next_consumed);  
    ...  
    <data in next consumed>  
}
```

Named pipes

- Accessed as files by processes
 - No parent child relation is necessary
 - Still behave like FIFO queues (even called fifo)
 - Several processes can use the named pipe

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```
char myfifo = '/tmp/myfifo';  
mkfifo (myfifo, 0666); // creates the fifo or named pipe  
  
...  
fd = open(myfifo, O_WRONLY); // Process A  
write(fd, ...);              // Process A  
close(fd);                    // Process A  
  
...  
fd = open(myfifo, O_RDONLY); // Process B  
read(fd, ...);               // Process B  
close(fd);                    // Process B
```

